## **Algorithms**

H. J. WEGSTEIN, Editor

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ALGORITHM 77
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INTERPOLATION, DIFFERENTIATION, AND IN-TEGRATION

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real procedure AVINT (nop, jt, xarg, xlo, xup, xa, ya); value nop, jt, xarg, xlo, xup; real xarg, xlo, xup; integer nop, it; real array xa, ya;

comment This procedure will perform interpolation, differentiation, or integration operating upon functions of one variable which over part or all of the interval of interest are adequately described by a di-parabolic fit.

The routine was originally programmed as an open subroutine for the IBM 704 in FORTRAN II and occupied 323 memory locations. It is based upon a Lagrange interpolation scheme specialized for averaged second order parabolas. The technique finds the slope of a function numerically defined at points 1, 2, 3 and 4 by fitting a parabola through the points 1, 2, 3, and another parabola through the points 2, 3, and 4. The slope then, at point 2, is the average analytical derivative of the two parabolas, i.e. the coefficients of the parabola through points 1, 2 and 3  $(a_1x_2^2+b_1x_2+c_1)$  and the coefficients of the parabola through points 2, 3, and 4  $(a_2x_2^2+b_2x_2+c_2)$ are determined by applying Lagrange's equations as shown below. The arithmetic mean of these coefficients  $a = (a_1 + a_2)/2$ ,  $b = (b_1+b_2)/2$ ,  $c = (c_1+c_2)/2$  are used to supply the slope in the interval from 2 to 3, namely (2ax + b).

The interpolation is calculated in similar fashion, except the final formula is that a parabola  $(ax^2 + bx + c)$ .

The integration is performed likewise by a curve fitting process, e.g. the integral between any two points say 2 and 3 is the average integral of the two parabolas between the independent coordinate limits for points 2 and 3. The averaging process is done for each interval along the abscissa as the results obtained are accumulated to evaluate the definite integral.

Applying Lagrange's equations, the coefficients a, b, and c may be found by defining:  $T_i = y_i / \prod_{i=1, i \neq j}^n (X_i - X_i)$  where  $\begin{array}{l} y \, = \, f(x), \quad n \, = \, 3, \quad j \, = \, 1, \, 2, \, \cdots, \, n, \ then \ a \, = \, \sum_{i=1}^n \, T_i \, , \\ b \, = \, \sum_{i=1}^n \, T_i \sum_{j=1, \ j \neq i}^n \, X_j \, , \quad c \, = \, \sum_{i=1}^n \, T_i \prod_{j=1, \ j \neq i}^n \, X_j \, ; \end{array}$ 

begin real ca, cb, cc, a, b, c, syl, syu, term1, term2, term3, da,

integer

switch alpha := L1, L1, L12; switch beta := L9, start:

switch gamma := L10, L11; switch delta := L8,

For interpolation, differentiation or integration set comment

L1: if  $xarg \ge xa [nop-1]$  then go to L2; if  $xarg \le xa$  [1] then go to L3;

AVINT := da; go to exit; exit1: AVINT := dif; go to exit; exit2:

AVINT := sum;exit3:

end exit:

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dif. sum:
jm, js, jul, ia, ib;
  L5, L6;
  L8, L13;
  jt = 1, 2, or 3 respectively;
go to alpha [jt];
if xarg \ge xa [nop] then go to L2;
if xarg \le xa [2] then go to L3; go to L4;
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jm := nop-1; js := 1; go to term;
   L2:
   L3:
             jm := 2; js := 1; go to term;
comment
             Locate argument;
             for ia := 2 step 1 until nop do begin
   L4:
             if xa [ia] > xarg then go to L7; jm := ia end;
             Before loop is complete xarg \leq xa [ia];
comment
             ca := a; cb := b; cc := c; js := 3; jm :=
   L5:
              jm+1; go to term;
             a := (ca+a)/2; b := (cb+b)/2; c := (cc+c)/2;
   L6:
              go to L9;
             js := 2; go to term;
   L7:
             go to beta [js];
   L8:
    L9:
             go to gamma [jt];
comment
             Interpolation, jt = 1;
             da := a \times xarg \uparrow 2 + b \times xarg + c; go to exitl;
   L10:
             Differentiation, jt = 2;
comment
             dif := 2 \times xarg + b; go to exit2;
    L11:
             Integration, jt = 3;
comment
             sum := 0; syl := xlo; jul := nop - 1;
    L12:
               ib := 2;
             for jm := ib step 1 until jul do begin;
    L16:
comment
             Lagrange formulae;
             term1 := ya [jm - 1]/((xa [jm - 1] - xa[jm]) \times
               (xa[im - 1] - xa[im + 1]));
             term2 := ya [jm]/((xa [jm] - xa [jm - 1]) \times
                (xa[jm] - xa[jm + 1]));
             term3 := ya [jm + 1]/((xa [jm + 1] - xa [jm - 1]) \times
               (xa [jm + 1] - xa [jm]));
             a := term1 + term2 + term3;
             b := -(xa [jm] + xa [jm + 1]) \times term1 - (xa
               [jm - 1] + xa [jm + 1]) \times term2 - (xa [jm - 1] +
               xa [jm]) × term3;
             c := xa [jm] \times xa [jm + 1] \times term1 + xa [jm - 1] \times
                xa [jm + 1] \times term2 + xa [jm - 1] \times xa [jm] \times
                term3; go to delta [jt];
    L13:
             if jm \neq 2 then go to L14;
             ca := a; cb := b; cc := c; go to L15;
    L14:
             ca := (a + ca)/2; cb := (b + cb)/2; cc :=
                (c + cc)/2;
             syu := xa [jm];
    L15:
             sum := sum + ca × (syu \uparrow 3 - syl \uparrow 3)/3 + cb ×
                (\text{syu} \uparrow 2 - \text{syl} \uparrow 2)/2 + \text{cc} \times (\text{syu} - \text{syl});
             ca := a; cb := b; cc := c; syl := syu end;
comment
             End of loop on [jm] index;
             sum := sum + ca × (xup ↑ 3-sy1 ↑ 3)/3 + cb ×
                (\sup \uparrow 2\text{-sy1} \uparrow 2)/2 + \text{ce} \times (\sup - \text{sy1}); \text{ go}
                to exit3;
             ib := jm; jul := ib; go to L16;
    term:
             The results for interpolation, differentiation, and
comment
                integration are da, dif, and sum respectively;
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## RATIONAL ROOTS OF POLYNOMIALS WITH IN-TEGER COEFFICIENTS

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comment This Algol procedure, named ratfact, for finding rational roots of polynomials with integer coefficients is a pedagogical example illustrating the use of the for statement described in section 4.6.3. Also, an extension suggested by J. Peck of the well-known polynomial evaluation by nesting, i.e. Horner's method, is used. The polynomial  $f(x) = a_0 + a_1x +$  $\cdots + a_n x^n$  with integer coefficients and with  $a_0 a_n \neq 0$  has a lowest term rational root p/q if and only if  $a_0q^n + a_1q^{n-1}p +$  $\cdots$   $+a_{n-1}q$   $p^{n-1}$  +  $a_{n}p^{n}$  = 0, also q must be a factor of  $a_{n}$  and p a factor of an Procedure Rateact outputs the nonzero rational roots p/q by execution of the procedure whose formal name is print. The output procedure uses the string whose formal name is format for control of the output format;

procedure ratfact (a, n, print, format);

integer array a[0:n]; integer n; procedure print; string format;

begin integer i, p, q, r, t, f, g;

```
p loop: for p := 1 step 1 until abs (a[0]) do
```

begin comment if p is not a factor of a [0] or q is not a factor of a[n] then skip to the end of the loop for advance in the respective for list:

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if a[0] \neq (a[0] \div p) \times p then go to 1
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else q loop: for q := 1 step 1 until abs (a[n]) do begin if  $a[n] \neq (a[n] \div q) \times q$  then go to 2

begin comment root test and print;

comment start polynomial evaluation;

f := g := a[0]; t := p;

for i := 1 step 1 until n do

begin  $r := a[i] \times t$ ;

 $f := f \times q + r;$ 

 $g := -g \times q + r;$ 

 $t := t \times p;$ 

end polynomial evaluation;

comment computing r saves one subscript evaluation:

if f=0 then print (format, p, q);

if g=0 then print (format,-p, q);

comment print is the formal name of the procedure to be used to output the variables in the format specified by the string whose formal name is format; end root test and print;

2: end q loop;

1: end p loop:

end ratfact, without overflow test.

## ALGORITHM 79 DIFFERENCE EXPRESSION COEFFICIENTS

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procedure dicol (k, n, xp, xtab, coef);
value k, n; integer k, n; real xp;
array xtab, coef;
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comment dicol produces the coefficients for the n ordinates (corresponding to the abscissae, xtab) in the n-point finite difference expression for the k-th derivative evaluated at xp. The method used is to determine the analytic expression for the k-th derivative of each coefficient in the n-point Lagrangian interpolation formula and evaluate it at xp. Note that k=0 will produce the Lagrangian interpolation coefficients themselves;

begin integer array xuse [1 : n-1]; real factk, sum, denom, part:

integer i, terms, j, m, high;

factk := 1.0; for i := 2 step 1 until k do factk := i×factk; terms := n-k-1; if terms <0 then go to Z;

for j := 1 step 1 until n do

```
loop: begin sum := 0; denom := 1.0; part := 1.0;
           for i := 1 step 1 until n do
           if i \neq j then denom := denom\times (xtab [j] - xtab [i]);
           if terms = 0 then go to Y;
           m := 1; high := 1;
```

A: if  $(high = j) \lor (xtab \{high\} = xp)$  then A1: begin high := high + 1; go to A end A1; if high > n then A2: begin m := m-1; if m>0then

A3: begin high := xuse [m]+1; go to A end A3; go to X end A2;

xuse [m] := high; m := m+1;

if m≤terms then begin high := high + 1; go to

for i := 1 step 1 until terms do  $part := part \times (xp - xtab [xuse [i]]);$ sum := sum + part; m := terms; part := 1.0;high := xuse [terms] + 1; go to A;

Y: sum := 1.0;

X: coef [j] := sum × factk/denom end loop; go to EXIT;

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Z: for i := 1 step 1 until n do coef [i] := 0;

EXIT: end dicol

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