

# Algorithms

J. H. WEGSTEIN, Editor

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## ALGORITHM 129

### MINIFUN

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**procedure** *MINIFUN* (*t1*, *b1*, *eps*, *n*, *ncnt*, *fmin*, *xmin*, *k1*, *GFUN*);

**value** *t1*, *b1*, *eps*, *n*, *ncnt*; **integer** *n*, *ncnt*, *k1*; **real** *fmin*;  
**real procedure** *GFUN*; **array** *t1*, *b1*, *eps*, *xmin*;

**comment** *MINIFUN* is a subroutine to find the minimum of a function of *n* variables, using the method of steepest descent. Input is:

1. *t1*(*i*), *i* = 1, 2, ..., *n*, the upper limits of the search region
2. *b1*(*i*), *i* = 1, 2, ..., *n*, the lower limits of the search region
3. *eps*(*i*), *i* = 1, 2, ..., *n*, the convergence criteria. The function must be a minimum in the region  $|x(i) - xmin(i)| \leq eps(i)$
4. *n*, the number of variables (the dimension of the arrays)
5. *ncnt*, the maximum number of iterations. The routine searches for a minimum until  $|x(i) - xmin(i)| \leq eps(i)$  for all *i*, or until *icnt* = *ncnt*, whichever happens first.

Output is:

1. *fmin*, the minimum value of the function
2. *xmin*(*i*), *i* = 1, ..., *n*, the point at which the minimum occurs
3. *k1*, an error code

If *k1* = 1, a minimum has been found within the specified number of iterations and the minimum is less than all

values of the function at the centers of the planes forming the boundary of the epsilon-cube

If *k1* = 2,  $\Delta x(i) \leq eps(i)$  but a new minimum has been found

If *k1* = 3, *ncnt* has been exceeded without  $\Delta x(i) \leq eps(i)$ . In this case, a test is made to see if the current minimum is a minimum in the epsilon-cube.

*MINIFUN* has been written as a FORTRAN II subroutine and is available from the SMSA Computation Center. It should be noted that the FORTRAN II deck has been tested only on some relatively simple functions of two variables, such as *GFUN* (*x,y*) = cos(*xy*). The writer does not claim that the algorithm has been thoroughly tested;

```
begin integer j, i, icnt, k; real w, dmax, alamb, ft;  
  array wnew [1:n], xt[1:n], x1b[1:n], xub [1:n],  
    del1x[1:n], d12x[1:n], xmin[1:n], x[1:n, 1:4], g[1:n, 1:4],  
    dxmin[1:n], d2xmn[1:n];  
  comment start looking for a minimum at midpoint of region;  
  for j := 1 step 1 until n do  
    begin wnew[j] := (t1[j] + b1[j])/2; xt[j] := wnew[j];  
      xub[j] := t1[j]; x1b[j] := b1[j]; del1x[j] := (xub[j]  
        - x1b[j])/5;  
      d12x[j] := del1x[j]↑2; xmin[j] := xt[j]  
    end;  
  fmin := GFUN (xmin);  
  for j := 1 step 1 until n do  
    begin w := xt[j]; for i := 1 step 1 until 4 do  
      begin x[j, i] := x1b[j] + i × del1x[j];  
        xt[j] := x[j, i]; g[j, i] := GFUN (xt);  
      end;  
      xt[j] := w;  
      dxmin[j] := (g[j, 3] - g[j, 2])/del1x[j];  
      d2xmn[j] := (g[j, 4] - g[j, 3] - g[j, 2] + g[j, 1])/d12x[j]  
    end;  
  comment first and second difference quotients have been computed;  
  icnt := 0; dmax := dxmin[1]; k := 1;  
  nustep: for j := 2 step 1 until n do  
    begin if abs(dmax) < abs(dxmn[j]) then  
      begin dmax := dxmn[j]; k := j  
    end;  
  end;  
  alamb := dxmin[k]/d2xmn[k]; w := xt[k] - alamb;  
  comment a new coordinate has been computed for the variable having the largest first partial derivative. It will be checked to see if the new point still lies within the region and search will continue;  
  if w < b1[k] then w := b1[k] else if w > t1[k] then w := t1[k];  
  xt[k] := w; ft := GFUN (xt);  
  if ft < fmin then go to check else  
  restart: if xt[k] < wnew[k] then go to 1bdchk  
    else if xt[k] = wnew[k] then go to stnubds  
    else if t1[k] > xt[k] then go to nupbds  
    else xt[k] := 1.5 × wnew[k];  
  nupbds: xub[k] := t1[k]; x1b[k] := 2 × xt[k] - t1[k]; go to newdel;  
  stnubds: x1b[k] := xt[k] - 0.5 × wnew[k]; xub[k] := xt[k] + 0.5 × wnew[k];  
  newdel: del1x[k] := 0.2 × (xub[k] - x1b[k]); d12x[k] := del1x[k]↑2;  
  for i := 1 step 1 until 4 do  
    begin x[k, i] := x1b[k] + i × del1x[k]; w := xt[k];  
      xt[k] := x[k, i]; g[k, i] := GFUN (xt); xt[k] := w  
    end;  
  dxmin[k] := (g[k, 3] - g[k, 2])/del1x[k];  
  d2xmn[k] := (g[k, 4] - g[k, 3] - g[k, 2] + g[k, 1])/d12x[k];  
  icnt := icnt + 1;  
  if icnt > ncnt then go to outcd else go to nustep;  
  1bdchk: if xt[k] ≤ b1[k] then xt[k] := 0.5 × wnew[k]
```

```

    else  $x1b[k] := b1[k]$ ;  $xub[k] := 2.0 \times xt[k] - b1[k]$ ;
  go to newdel;
check:  $fmin := ft$ ;  $xmin[k] := xt[k]$ ;
  for  $j := 1$  step 1 until  $n$  do if  $delx[j] > eps[j]$  then go to restart;
recheck: for  $j := 1$  step 1 until  $n$  do
  begin  $w := xmin[j]$ ;  $xmin[j] := w + eps[j]$ ;  $ft := GFUN(xmin)$ ;
    if  $ft < fmin$  then go to set2;  $xmin[j] := w - eps[j]$ ;
     $ft := GFUN(xmin)$ ; if  $ft < fmin$  then go to set2;  $xmin[j] := w$ 
  end;
if  $k1 < 3$  then  $k1 := 1$ ; go to bgend;
set2:  $k1 := 2$ ; go to bgend;
outcd:  $k1 := 3$ ; go to recheck;
bgend: end MINIFUN;

```

### ALGORITHM 130

#### PERMUTE

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**procedure** PERMUTE ( $A, n, x$ )

**array**  $A$ ; **integer**  $n, x$ ;

**comment** Each entry into PERMUTE generates the next permutation of the first  $n$  elements of  $A$ . If  $A$  is read as a number ( $A[1]A[2] \dots A[n]$ ), each generation is larger than the last:  $n := 4, x := 1$

$A[1]$	1	1	1	8	8	8	
$A[2]$	1	8	8	1	1	8	
$A[3]$	8	1	8	1	8	1	Permutations = $\frac{4!}{2!2!}$
$A[4]$	8	8	1	8	1	1	end

Identical elements in  $A$  reduce the number of permutations. The array should be ordered before the first call on PERMUTE. Integer  $x$  specifies the first elements whose order should be preserved:  $n := 4, x := 3$

$A[1]$	1	1	1	4	
$A[2]$	2	2	4	1	
$A[3]$	3	4	2	2	Permutations = $\frac{4!}{3!}$
$A[4]$	4	3	3	3	end

Before the first call on PERMUTE for a given array, *first* should be made true. If *more* is true, then PERMUTE was able to give another permutation;

**begin array**  $B[1:n]$ ; **integer**  $f, i, k, m, p$ ; **real**  $r$ ; **own real**  $t$ ;

if *first* then  $t := A[x]$ ; *first* := false;

for  $i := 1$  step 1 until  $n$  do  $B[i] := 0$ ;

for  $i := n$  step -1 until 2 do

begin if  $A[i] > t \wedge A[i] > A[i - 1]$  then go to find; end;

*more* := false; go to exit;

find: for  $k := n$  step -1 until  $i$  do

begin if  $A[k] > t \wedge A[k] > A[i - 1]$  then

begin  $B[k] := A[k]$ ;  $m := k$ ; end; end;

for  $k := n$  step -1 until  $i$  do

begin if  $B[k] > 0 \wedge B[k] < B[m]$  then

begin  $B[m] := B[k]$ ;  $f := k$ ; end; end;

$r := A[i - 1]$ ;  $A[i - 1] := B[m]$ ;  $A[f] := r$ ;

schell:  $p := i - 1$ ;  $m := n - p$ ;

for  $m := m/2 - .4$  while  $m > 0$  do

begin  $k := n - m$ ;

for  $f := p + 1$  step 1 until  $k$  do

begin  $i := f$ ;

if  $A[i] > A[i + m]$  then  
begin  $r := A[i + m]$ ;  $A[i + m] := A[i]$ ;

$A[i] := r$ ;  $i := i - m$ ;

if  $i \geq p + 1$  then go to comp;

end end end schell;

comp: if  $A[i] > A[i + m]$  then

begin  $r := A[i + m]$ ;  $A[i + m] := A[i]$ ;

$A[i] := r$ ;  $i := i - m$ ;

if  $i \geq p + 1$  then go to comp;

end end end schell;

exit: end PERMUTE

### ALGORITHM 131

#### COEFFICIENT DETERMINATION\*

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\* This procedure pertains to research work sponsored in part by NSF Grant G-7361.

**procedure** DET ( $n, G, H$ );

**array**  $G, H$ ; **integer**  $n$ ;

**comment** Given the first  $n$  coefficients of the power series  $G(z) = g_1 + g_2z + g_3z^2 + \dots + g_nz^{n-1} + \dots$ , and  $H(z) = h_1 + h_2z + h_3z^2 + \dots + h_nz^{n-1} + \dots$ , this procedure determines the coefficients  $d_i, i = 1, \dots, n$ , of the power series which is the expansion of the quotient  $H(z)/G(z)$ . It is assumed that  $g_1 \neq 0$ . The arrays  $G$  and  $H$  initially contain the coefficients of  $G(z)$  and  $H(z)$ , respectively. The integer  $n$  is the number of known coefficients in the expansion of  $G(z)$  and  $H(z)$ . At the conclusion,  $H_i$  contains the coefficient  $d_i$ . The procedure may also be useful in calculating residues for certain complex functions. Suppose  $F(z) = H(z)/G(z)$  is a complex valued function of a complex variable and that  $F$  has a pole of order  $m$  at  $z = b$ , where  $H(z) = \sum_{k=1}^m h_k(z - b)^{k-1}$ ,  $G(z) = \sum_{k=1}^{\infty} g_k(z - b)^{k+m-1}$ , and  $g_1 \neq 0, h_1 \neq 0$ . The required residue at  $z = b$  is  $d_m$  where

$$D(z) = \left[ \sum_{k=1}^{\infty} h_k(z - b)^{k-1} \right] / \left[ \sum_{k=1}^{\infty} g_k(z - b)^{k-1} \right]$$

$$= \sum_{j=1}^{\infty} d_j(z - b)^{j-1}.$$

For more on this, one is referred to Einar Hille, "Analytic Function Theory, Vol. I," Ginn and Co., 1959, pages 242-244;

**begin integer**  $i, j, n$ ; **real**  $alpha, beta$ ;

$alpha := 1/G[1]$ ;

for  $j := 1$  step 1 until  $n$  do

begin  $beta := alpha \times H[j]$ ;

for  $i := j + 1$  step 1 until  $n$  do

$H[i] := H[i] - (beta \times G[i - j + 1])$  end;

for  $j := 1$  step 1 until  $n$  do

$H[j] := H[j] \times alpha$ ;

end

DET

### ALGORITHM 132

#### QUANTUM MECHANICAL INTEGRALS OVER ALL SLATER-TYPE INTEGRALS

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**real procedure:** *allslater* ( $p, q, pe, qe, np, nq, lp, lq, mp, mq, na, nb$ )

internuclear distance: ( $r$ );

**real**  $pe, qe, r$ ; **integer**  $p, q, np, nq, lp, lq, mp, mq, na, nb$ ;

**comment** The Slater-type orbitals frequently used in quantum mechanical calculations on atoms and molecules are defined as  $p = k(np, pe) r^{np-1} e^{-(pe)r} Y_{lp}^{mp}(\theta, \phi)$ , where  $k(np, pe)$  is a normalization constant,  $Y_{lp}^{mp}(\theta, \phi)$  is a spherical harmonic with the phase convention  $[Y_{lp}^{mp}(\theta, \phi)]^* = (-1)^m Y_{lp}^{-m}(\theta, \phi)$ ,  $np$  is a positive integer,  $lp$  is an integer,  $lp < np$ ,  $mp$  is an integer,  $-lp \leq mp \leq lp$ ; and  $pe$  is a real positive constant. Algorithm 110, Y. A. Kruglyak and D. R. Whitman (*Comm. ACM*, July 1962) serves to compute integrals over certain operators of a quite restricted class of Slater-type orbitals,  $np \geq 4, lp = 1, mp = 0$ . The algorithm given here will compute all integrals of the form

$$\int p_c(r_c^c) q_c dr$$

which can be expressed in terms of the simple  $A_n(b)$  and  $B_n(a)$  functions. The subscript  $c$  denotes either of the two nuclei of

a diatomic molecule. These integrals include all those one-electron integrals necessary for a conventional energy calculation on a diatomic molecule. In the arguments of *allslater* *p* and *q* are numerical designations for the respective orbitals. *p* and *q* are even or odd as they respectively are associated with the "left," *a*, nucleus or "right," *b*, nucleus of a diatomic molecule. Global arrays, *fact* 1, of factorials and *binom*, of binomial coefficients are assumed. We first define some procedures utilized by *allslater*. The main program begins at the label *set*;

```

begin real norm, r2, alpha, beta, s, clp, clq, bpci;
integer nsum, lsum, peven, qeven, podd, godd, limitp, limitq,
g, h, i, j, n1p, n1q, lmp, lmq, gama, gamb, aidaa, aidab, gam,
aida, num2; real array avalues [0:21], bvalues[0:21]; real procedure
c1, bpc, modulus;
real procedure c1(l,m,j); value l,m,j, integer l,m,j;
begin c1 := ((-1)j × fact1[2×(l - j)] / ((2j) × fact1
[l-2×j-m] ×
fact1[l - j] × fact1[j])
end c1;
real procedure bpc(i, j, k); value i,j,k, integer i,j,k;
begin real t; integer m; t := 0;
for m := 0 step 1 until k do
begin t := t + ((-1)k - m
× binom [i, m] × binom [j, k - m]
end
end bpc;
real procedure modulus (i, j); value i, j; integer i, j;
begin modulus := 1 - abs(i ÷ j) × j
end modulus;
procedure avector (b, nmax, avalues); value b, nmax;
real b; integer nmax; real array avalues;
integer m;
avales[0] = exp(-b)/b;
if nmax = 0 then go to exit;
for m = 1 step 1 until nmax do
begin avalues[m] = avalues[0] + (m/b) × avalues[m - 1]
end;
exit: end avector;
procedure bvector(a nmax, bvalues); value a, nmax; real a;
integer nmax; real array bvalues; real procedure modulus;
comment This procedure computes a sequence of values for the
integral,  $B_n(a) = \int_{-1}^1 x^n e^{-ax} dx$ , for  $n = 0$  to  $n = nmax$ . If  $a \geq alim$ 
then  $B_0(a)$  is computed and upward recursion is used to generate the higher  $n$ 
values. If  $a < alim$  then  $B_{nmax}(a)$  is computed by series expansion and downward
recursion is used to generate the smaller  $n$  values. alim is determined within the
program by a simplification of a result of Gautschi (J. ACM 8, 21 (1961)).
Gautschi has made an analysis of the recursive procedures for the  $B_n(a)$  which
could be taken as a model for workers in molecular quantum mechanics;
begin real fxx, fxy, numerator, denom, sum, factor1, tsum
factor2, t, aa; integer m, mn;
begin if abs(a)  $\geq ((nmax+nmax/6+3)/2.3)$  then
up: begin fxx := exp(a);
fxy := 1/fxx;
bvalues [0] := (fxx-fxy)/a;
for m := 1 step 1 until nmax do
begin fxx := - fxx;
bvalues[m] := (fxx-fxy + m ×
bvalues[m-1])/a
end;
go to exit;
end up;
down: begin aa := axa;
if modulus (nmax, 2)  $\neq$  0 then
setodd: begin numerator := nmax + 2;
sum := a/numerator;
factor1 := -2;

```

```

factor2 := 3;
go to compute;
end setodd;
seteven: begin numerator := nmax + 1;
sum := 1/numerator;
factor1 := factor2 := 2;
end seteven;
compute: begin denom := numerator + 2;
t := sum;
t := (((t/factor2) × aa)
/(factor2-1)) × numerator)
/denom;
tsum := t + sum;
if (sum-tsum) = 0 then
begin bvalues[nmax] := sum × factor1;
go to recur;
end;
begin factor2 := factor2 + 2;
numerator := denom;
sum := tsum;
go to compute;
end compute;
recur: begin fxx := exp(a);
fxy := 1/fxx;
mn := nmax - 1;
if modulus(nmax, 2)  $\neq$  0 then
fxx := -fxx;
for m := mn step -1 until 0 do
begin fxx = -fxx;
bvalues[m] := (fxx+fxy + a ×
bvalues[m+1])/(m+1);
end
end recur;
end down;
end;
exit: end bvector;
set: begin if (mp + mq)  $\neq$  0 then
begin allslater := 0.0; go to exit end;
set: begin norm := sqrt (((2×pe)j
(2×np+1) × (2×lp+1) × fact1[lp-mp] × (2×qe)j
(2×nq+1) × (2×lq+1) × fact1[lq-mq]) / (fact1[2×
np] × fact1[lp+mp] × fact1[2×nq] × fact1[lq+mq] ×
4));
nsum := np+nq;
lsum := lp+lq;
r2 := r/2;
norm := norm × (r2j(nsum+1+na+nb));
alpha := r2 × (pe+qe);
beta := r2 × (((-1)jp) × pe + ((-1)jq) × qe);
num2 := 2;
avector (alpha, nsum, avalues);
bvector (beta, nsum, bvalues);
peven := modulus (p+1,2);
qeven := modulus (q+1,2);
podd := modulus (p,2);
qodd := modulus (q,2);
limitp := (lp-mp) ÷ num2;
limitq := (lq-mq) ÷ num2;
s := 0;
end set;
sum: begin for g := 0 step 1 until limitp do
begin c1p := c1(lp,mp,g);
for h := 0 step 1 until limitq do
begin c1q := c1(lq,mq,h);
n1p := np-lp+2×g-1;
n1q := nq-lq+2×h-1;
lmp := lp-mp-2×g;
lmp := lq-mq-2×h;

```

```

gama := n1p × peven + n1q × qeven + 1 + na;
gamb := n1p × podd + n1q × godd + 1 + nb;
aidaa := lmp × peven + lmq × qeven;
aidab := lmp × podd + lmq × godd;
gam = gama + gamb;
aida = aidaa + aidab;
for i := 0 step 1 until gam do
begin bpci := bpc(gama, gamb, i);
for j := 0 step 1 until aida do
begin
s := s + c1p × c1q × bpci ×
bpc(aidaa, aidab, j) ×
avalues[nsum+na+nb-i-j]
× bvalues[lsum-2 × (g+h) + i-j];
end
end
end;
allslater := s × norm;
end sum;
exit: end;
end allslater;

```

### ALGORITHM 133

#### RANDOM

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**real procedure** *RANDOM* (*A*, *B*, *X0*);

**value** *A*, *B*, *X0*;

**real** *A*, *B*;

**integer** *X0*;

**comment** *RANDOM* generates a rectangular distributed pseudo-random number in the interval  $A < B$ . *X0* is an integer starting-value. The first time *RANDOM* is used in a program *X0* should be a positive odd integer with 11 digits,  $X0 < 2^{36} = 34\,359\,738\,368$ . The following times *RANDOM* is used, *X0* should be  $X0 = 0$ . The mathematical method used is  $X_{n+1} = 5 X_n \pmod{2^{35}}$ . This sequence has period  $2^{35}$ . *RANDOM* was successfully run on *FACIT EDB* using *FACIT-ALGOL 1*, which is a realization of *ALGOL 60* for *FACIT EDB*, except for the declarator **own**, which is not included in *FACIT-ALGOL 1*. To test *RANDOM*, we computed  $1/N \sum X_n$  and  $1/N \sum X_n^2$  in the interval 0,1 for  $N = 500, 1000, 5000$ . The starting-value was  $X0 = 28\,395\,423\,107$ . The results were 0.50625, 0.48632, 0.50304 and 0.34304, 0.31681, 0.33469. Theoretically one expects 0.50000 and 0.33333;

**begin**

**integer** *M35*, *M36*, *M37*;

**own integer** *X*;

**if**  $X0 \neq 0$  **then begin**

*X* := *X0*; *M35* := 34 359 738 368; *M36* := 68 719 476 736;

*M37* := 137 438 953 472 **end**; *X* := 5 × *X*;

**if**  $X \geq M37$  **then** *X* := *X* - *M37*;

**if**  $X \geq M36$  **then** *X* := *X* - *M36*;

**if**  $X \geq M35$  **then** *X* := *X* - *M35*;

*RANDOM* := *X*/*M35* × (*B* - *A*) + *A* **end**

### ALGORITHM 134

#### EXPONENTIATION OF SERIES

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**procedure** *SERIESPWR*(*A*, *B*, *P*, *N*);

**comment** This procedure calculates the coefficients  $B[i]$  for the series  $(f(x))^P \equiv g(x) \doteq 1 + \sum B[i] \times x \uparrow i$ , ( $i = 1, 2, \dots, N$ ) given the coefficients of the series  $f(x) = 1 + \sum A[i] \times x \uparrow i$ . *P* may be any real number;

**value** *A*, *P*, *N*;

**array** *A*, *B*;

**integer** *N*;

**begin integer** *i*, *k*;

**real** *p*, *s*;

*B*[1] := *P* × *A*[1];

**for** *i* := 2 **step 1 until** *N* **do**

**begin** *s* := 0;

**for** *k* := 1 **step 1 until** *i*-1 **do**

*S* := *s* + (*P* × [*i*-*k*] - *k*) × *B*[*k*] × *A*[*i*-*k*];

*B*[*i*] := *P* × *A*[*i*] + (*s*/*i*)

**end for** *i*;

**end** *SERIESPWR*

### ALGORITHM 135

#### CROUT WITH EQUILIBRATION AND ITERATION

WILLIAM MARSHALL McKEEMAN\*

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\* This work was supported in part by the Office of Naval Research under contract Nonr 225(37).

**procedure** *LINEAR SYSTEM* (*A*) order:(*n*) right-hand sides:(*B*) number of right-hand sides:(*m*) answers:(*X*) determinant:(*det*, *ex*) condition of *A*:(*cnr*);

**integer** *n*, *m*, *ex*; **real array** *A*, *B*, *X*;

**comment**, *LINEAR SYSTEM* uses Crout's method with row equilibration, row interchanges and iterative improvement for solving the matrix equation  $AX = B$  where *A* is  $n \times n$  and *X* and *B* are  $n \times m$ . As special cases one sees that: for  $m \leq 0$ , only the determinant of *A* is evaluated, for  $m = 1$ , the algorithm solves a system of *n* equations in *n* unknowns, for  $m = n$  and *B* = the identity matrix, the algorithm inverts *A*.

If the algorithm breaks down for a singular or nearly singular matrix *A*, exit to a non-local label "singular" is provided. Five auxiliary procedures: *EQUILIBRATE*, *CROUT*, *PRODUCT*, *RESIDUALS* and *SOLVE* are declared with appropriate comments after the end of this procedure. This code is the result of the joint efforts of G. Guthrie, W. McKeeman, Cleve Moler, Margaret Salmon, Alan Shaw and R. Van Wyk. It was written following ideas presented by J. H. Wilkinson as a visiting lecturer in Professor George E. Forsythe's class in Advanced Numerical Analysis at Stanford, 1962;

**begin integer array** *pivot* [1:*n*]; **integer** *i*, *j*, *k*; **real** *mx*;

**real array** *LU*[1:*n*, 1:*n*], *y*, *res*, *mult*[1:*n*];

**comment**, remove appropriate factors from the rows of *A*... ; *EQUILIBRATE*(*A*, *n*, *mult*);

**comment** ... and save the result for the eventual computation of residuals during iteration;

**for** *i* := 1 **step 1 until** *n* **do**

**for** *j* := 1 **step 1 until** *n* **do** *LU*[*i*,*j*] := *A*[*i*,*j*];

**comment**, decompose the matrix into triangular factors;

*CROUT*(*LU*, *n*, *pivot*, *det*);

**comment**, assuming that there was no exit to "singular", evaluate the determinant in the form  $det \times (10.0 \uparrow ex)$ ;

**for** *i* := 1 **step 1 until** *n* **do** *y*[*i*] := *LU*[*i*,*i*] × *mult*[*i*];

*det* := *det* × *PRODUCT*(*y*, 1, *n*, *ex*);

**comment**, now begin to process right-hand sides;

**for** *k* := 1 **step 1 until** *m* **do**

**begin integer** *i*, *count*, *limit*; **real** *normy*, *kr*;

*kr* := *k*;

**comment**, scale the right-hand side;

**for** *i* := 1 **step 1 until** *n* **do** *res*[*i*] := *B*[*i*,*k*] := *B*[*i*,*k*]/*mult*[*i*];

**comment**, store the first approximation and its *L*(1) *norm*;

*normy* := 0;

```

SOLVE(LU, n, res, pivot, y);
for i := 1 step 1 until n do
begin
  normy := normy + abs(y[i]);
  X[i,k] := y[i]
end;
comment, enter the iterating loop. The iteration is terminated on the integer "limit" which itself is determined on the basis of the success of the first iteration and a machine-dependent real number designated here by "eps". For "eps", the programmer must insert the largest real number such that  $eps + 1.0 = 1.0$ ;
for count := 1, 2 step 1 until limit do
begin integer i; real t;
  comment, compute the residuals of the solution y;
  RESIDUALS(A,n,B,k,X,res);
  comment ... and find the next increment to the solution;
  SOLVE(LU,n,res,pivot,y);
  comment, set up termination conditions;
  if count = 1 then
  begin real normdy;
    normdy := 0;
    for i := 1 step 1 until n do normdy := normdy+abs(y[i]);
    if normdy = 0 then begin cnr := 1.0; go to enditer end;
    t := normy/normdy;
    comment, The quantity  $\|A\| \cdot \|A^{-1}\|$  (spectral norm) is called the condition number of the matrix A. It is a measure of the difficulty in solving the input equation and appears naturally in error bounds for the solution (see Wilkinson [3]). cnr is a direct measure of the error and experimentally approximates the condition number;
    cnr := ((kr - 1.0) × cnr + 1.0/(eps × t))/kr;
    if t < 2.0 then go to singular;
    limit := ln(eps)/ln(1.0/t);
  end;
  comment, store the new approximation;
  for i := 1 step 1 until n do X[i,k] := X[i,k] + y[i];
end iteration;
enditer:
end right-hand sides
end LINEAR SYSTEM;
procedure EQUILIBRATE (A) order:(n) multipliers:(mult);
integer n; real array A, mult;
comment, scaling the rows of the matrix A to roughly the same maximum magnitude (here, dividing by the largest element) allows the procedure CROUT to select effective pivotal elements for the Gaussian decomposition of the matrix. The iterating procedure will converge to the solution for the equilibrated matrix rather than the input matrix. If the matrix is badly conditioned then the solution is sensitive to perturbations in the input and the scaling division must be done not by the largest element but rather by the power of the machine number base (2 and 10 for binary and decimal machines, respectively) nearest the largest element so as to avoid rounding errors. Equilibration is discussed in reference [3] p. 284;
begin integer i; real mx;
  for i := 1 step 1 until n do
  begin integer j;
    mx := 0.0; comment, find the largest element;
    for j := 1 step 1 until n do
      if abs(A[i,k]) > mx then mx := abs(A[i,k]);
    if mx = 0.0 then go to singular;
    comment, now store the multiplier and scale the row;
    mult[i] := mx; comment := base ↑ ex for exact scaling;
    if mx ≠ 1.0 then
      for j := 1 step 1 until n do A[i,j] := A[i,j]/mx
    end
  end
end EQUILIBRATE;

```

```

procedure CROUT (A) order:(n) pivots:(pivot) interchanges:(sg).
integer n; integer array pivot; real array A; real sg;
comment, this is Crout's method with row interchanges as formulated in reference [1] for transforming the matrix A into the triangular decomposition LU with all the  $L[k,k] = 1.0$ . pivot[k] stores the index of the pivotal row at the k-th stage of the elimination for use in the procedure SOLVE;
begin integer i, j, k, imax, p; real t, quot;
  real procedure IP1 (A) extra term:(t) length:(f);
  integer f; real t; real array A; comment non-local i, j, k;
  comment, IP1 forms a row by column inner product of A, namely the sum of  $A[i,p] \times A[p,k]$  for  $p := 1, 2, \dots, f$ , and then adds the extra term t. If  $f < 1$ , the value of IP1 is t. This procedure is the inner loop of the algorithm. The programmer can expect a substantial advantage from substituting a faster and more accurate inner product here;
  begin real sum; integer p;
    sum := t;
    for p := 1 step 1 until f do sum := sum + A[i,p] × A[p,k];
    IP1 := sum
  end IP1;
  sg := 1.0;
  comment, k is the stage of the elimination;
  for k := 1 step 1 until n do
  begin
    t := 0;
    for i := k step 1 until n do
    begin comment, compute L. Note that the first calls on IP1 are empty;
      A[i,k] := -IP1(A, -A[i,k], k-1);
      if abs(A[i,k]) > t then
        begin t := abs(A[i,k]); imax := i end
    end;
    if t = 0 then go to singular;
    comment, A[imax,k] is the largest element in the remainder of column k. Interchange rows if necessary and record the change;
    pivot[k] := imax;
    if imax ≠ k then
    begin
      sg := -sg;
      for j := 1 step 1 until n do
      begin
        t := A[k,j]; A[k,j] := A[imax, j]; A[imax, j] := t
      end
    end;
    comment, compute a column of multipliers;
    quot := 1.0/A[k,k];
    for i := k+1 step 1 until n do A[i,k] := A[i,k] × quot;
    comment, and compute a row of U;
    for j := k+1 step 1 until n do
      A[k,j] := -IP1(A, -A[k,j], k-1)
    end
  end
end CROUT;
real procedure PRODUCT (factors) start:(s) finish:(f)
  exponent:(ex);
integer s,f,ex; real array factors;
comment, PRODUCT multiplies the numbers stored from index s through f inclusive in the array "factors", preventing exponent overflow. The answer is normalized so that  $1.0 > abs(PRODUCT) \geq 0.1$ . The exponent appears in ex;
begin integer i; real p, p1;
  ex := 0; p := 1.0;
  for i := s step 1 until f do
  begin
    p1 := factors [i];
    if abs(p1) < 0.1 then begin p1 = 10.0 × p1; ex := ex-1
    end;
    p := p × p1;
  end
end

```

```

if  $p = 0$  then begin  $ex := 0$ ; go to  $fin$  end;
1: if  $abs(p) < 0.1$  then
  begin  $p := p \times 10.0$ ;  $ex := ex - 1$ ; go to 1 end;
2: if  $abs(p) \geq 1.0$  then
  begin  $p := p/10.0$ ;  $ex := ex + 1$ ; go to 2 end;
end;
 $fin$ :  $PRODUCT := p$ 
end  $PRODUCT$ ;
procedure  $RESIDUALS$  ( $A$ )  $order:(n)$   $right-hand\ sides:(B)$ 
 $column\ of\ B:(k)$   $approximate\ solution:(X)$   $residuals:(res)$ ;
integer  $n, k$ ; real array  $A, B, X, res$ ;
comment,  $RESIDUALS$  computes  $b - Ay$  where  $b$  is the  $k$ th
 $column\ of\ the\ right-hand\ side\ matrix\ B$  and  $y$  is the  $k$ th  $column\ of\ X$ ;
  real procedure  $IP2$  ( $A$ )  $row:(i)$   $order:(n)$   $approximate$ 
 $solution:(X)$ 
 $column:(k)$   $extra\ term:(t)$ ;
  integer  $i, k, n$ ; real t real array  $A, X$ ;
  comment,  $IP2$  forms the inner product of row  $i$  of the matrix
 $A$  and column  $k$  of the solution matrix  $X$ , then adds the
  single term  $t$ . It is essential that  $IP2$  be an "accumulating"
  or double precision inner product as discussed in reference
  [3] p. 296. The value of  $IP2$  is the rounded single precision
  result of the double precision arithmetic. The body of the
  procedure is left undefined;
begin integer  $i$ ;
for  $i := 1$  step 1 until  $n$  do
   $res[i] := -IP2(A, i, n, X, k, -B[i, k])$ 
end  $RESIDUALS$ ;
procedure  $SOLVE$  ( $A$ )  $order:(n)$   $right-hand\ side:(b)$   $pivots:$ 
 $(pivot)$   $answer:(y)$ ;
integer  $n$ ; integer array  $pivot$ ; real array  $A, b, y$ ;
comment,  $SOLVE$  processes a right-hand side  $b$  and then back-
  solves for the solution  $y$  using the  $LU$  decomposition provided
  by  $CROUT$ ;
begin integer  $k, p$ ; real t;
for  $k := 1$  step 1 until  $n$  do
  begin
     $t := b[pivot[k]]$ ;  $b[pivot[k]] := b[k]$ ;
    for  $p := 1$  step 1 until  $k-1$  do  $t := t - A[k, p] \times b[p]$ ;
     $b[k] := t$ 
  end ...having modified  $b$  by  $L$  inverse;
  comment, now the back solution for  $y$ ;
  for  $k := n$  step -1 until 1 do
    begin
       $t := b[k]$ ;
      for  $p := k+1$  step 1 until  $n$  do  $t := t - A[k, p] \times y[p]$ ;
       $y[k] := t$ 
    end  $backsolution$ 
  end  $SOLVE$ 

```

#### REFERENCES

1. GEORGE E. FORSYTHE, Crout with Pivoting. Algorithm 16. *Comm. ACM* 3, 2 (Sept. 1960), 507.
2. DEREK JOHANN ROEK, Simultaneous System of Equations and Matrix Inversion Routine. Algorithm 92. *Comm. ACM* 5, 5 (May 1962), 286.
3. J. H. WILKINSON, Error Analysis of Direct Methods of Matrix Inversion, *J. ACM* 8, 3 (July 1961), 281-330.

#### ALGORITHM 136 ENLARGEMENT OF A GROUP M. WELLS\*

University of Leeds, England

\* Currently with Burroughs Corporation, Pasadena, California

```

procedure  $Enlarge\ group$  ( $G, n, g, Abelian$ );
array  $G, g$ ; integer  $n$ ; Boolean  $Abelian$ ;

```

**comment** This procedure combines the element  $g$  with the sub-group  $G$ , of  $n$  elements, to form a new group. The **Boolean**  $Abelian$  has the value **true** if the group to which  $G$  and  $g$  belong is Abelian. Two **procedures**,  $multiply$  and  $equal$  are assumed to be declared:  $multiply(G[i])$  by  $:(G[j])$  to give  $:(G[k])$  will set the element  $G_k$  equal to the product of the elements  $G_i$  and  $G_j$ .  $equal(G[i], G[j])$  is a **Boolean procedure** whose value is **true** if, and only if, the elements  $G_i$  and  $G_j$  are equal. On leaving the **procedure** the enlarged group is in  $G$ , and  $n$  is equal to the number of elements in the new sub-group  $G$ . The **procedure** will function correctly if  $g$  is included in  $G$  on entry. It is probable that  $g$  and the elements of  $G$  will be **arrays**, and the procedure body will, in practice, need to be altered considerably. The **procedure** has been used successfully in connection with problems of space-group theory;

```

begin integer  $i, j, k$ ;
for  $i := 1$  step 1 until  $n$  do
  if  $equal(G[i], g)$  then go to  $not\ new\ generator$ ;
   $n := n + 1$ ;  $G[n] := g$ ;
  for  $i := n$  step 1 until  $n$  do
    begin for  $j := 1$  step 1 until  $n$  do
      begin multiply  $(G[i], G[j], G[n+1])$ ;
        for  $k := 1$  step 1 until  $n$  do
          if  $equal(G[k], G[n+1])$  then go to  $not\ new\ element\ 1$ ;
           $n := n + 1$ ;
        not new element 1: if  $Abelian$  then go to  $take\ next\ element$ ;
         $multiply(G[j], G[i], G[n+1])$ ;
        for  $k := 1$  step 1 until  $n$  do
          if  $equal(G[k], G[n+1])$  then go to  $not\ new\ element\ 2$ ;
           $n := n + 1$ ;
        not new element 2: take next element;
      end of j-loop;
    end of i-loop;
  not new generator: end of group enlargement

```

#### ALGORITHM 137 NESTING OF FOR STATEMENT I DAVID M. DAHM & M. WELLS\* Burroughs Corp., Pasadena, Calif.

\* On leave of absence from the University of Leeds, England.

```

procedure  $Fors\ 1$  ( $n, P$ );
value  $n$ ; integer  $n$ ; procedure  $P$ ;
comment  $Fors\ 1$  generates a nest of  $n$  for statements with the
 $procedure\ P$  at their center. Two non-local arrays  $I$  and  $U$ ,
which give the value of the controlled variable and its upper
bound for each level are assumed to be declared;
begin integer  $j$ ;
if  $n = 0$  then  $P$ 
else for  $j := 1$  step 1 until  $U[n]$  do
  begin  $I[n] := j$ ;  $Fors\ 1(n-1, P)$  end end  $Fors\ 1$ 

```

#### ALGORITHM 138 NESTING OF FOR STATEMENT II DAVID M. DAHM & M. WELLS\* Burroughs Corp., Pasadena, Calif.

\* On leave of absence from the University of Leeds, England.

```

procedure  $Fors\ 2$  ( $P$ );
procedure  $P$ ;
comment  $Fors\ 2$  performs the same function as  $Fors\ 1$ , but is
more economic of storage space. It is expected, however,
that  $Fors\ 1$  would be more economic of time. The formal
parameter  $n$  is now replaced by the non-local integer  $n$ ;

```

```

begin if  $n = 0$  then  $P$ 
  else for  $I[n] := 1$  step 1 until  $U[n]$  do
    begin  $n := n - 1$ ; Fors 2 ( $P$ ) end;
   $n := n + 1$  end Fors 2

```

ALGORITHM 139  
SOLUTIONS OF THE DIOPHANTINE EQUATION  
J. E. L. PECK

University of Alberta, Calgary, Alberta, Canada

```

procedure Diophantus ( $a,b,c$ ); integer  $a,b,c$ ;
comment This procedure seeks the integer solutions of the
equation  $ax + by = c$ , where the integers  $a,b,c$  are given. It
assumes a non-local integer  $M$ , which should be as large as
storage will allow, two nonlocal labels INDETERMINATE
and NO SOLUTION and two non-local Boolean variables
'general solution' and 'time permits' which are self explanatory.
It also assumes the procedures abs, sign and print;
begin integer  $n,r,s,d,i$ ; integer array  $q[1:M]$ ;
 $n := i := 0$ ;  $d := s := \text{abs}(a)$ ;  $r := \text{abs}(b)$ ;
comment  $d$  will become the greatest common divisor of  $a$  and  $b$ .
  If  $b = 0$  then  $d = |a|$ . The vector  $q$  will retain the successive
  quotients in the Euclidean algorithm  $r_{i-1} = r_i q_i + r_{i+1}$ ,
   $i = 1, 2, \dots, n$ , where  $0 \leq r_{i+1} < r_i$ ,  $r_0 = |a|$ ,  $r_1 = |b|$ ,
  and  $r_{n+1} = 0$ ;
for  $i := i + 1$  while  $r \neq 0$  do
  begin  $n := i$ ;  $d := r$ ;  $q[i] := s \div d$ ;
   $r := s - d \times q[i]$ ;  $s := d$  end This records the quotients and
  the number  $n$  of divisions for use below;
if  $d = 0$  then go to if  $c = 0$  then INDETERMINATE
else NO SOLUTION; comment The case  $d = 0$  occurs when
 $a^2 + b^2 = 0$ . If  $d$  now does not divide  $c$  then the equation can-
not be solved so;
if  $(c \div d) \times d \neq c$  then go to NO SOLUTION;
if  $d \neq 1$  then
  begin  $a := a/d$ ,  $b := b/d$ ;  $c := c/d$  end, which removes
  the common factor and reduces the equation to the case
  where  $a$  and  $b$  are relatively prime;
  begin comment We shall now find  $u_1$  and  $v_1$  in order to
  express
   $1 = au_1 + bv_1$ , using the relations  $r_n = r_i v_i + r_{i-1} u_i$ ,
   $i = n, n-1, \dots, 1, v_n = 1, u_n = 0$ , and  $r_{i+1} = -r_i q_i + r_{i-1}$ ,
   $i = n-1, n-2, \dots, 1$ ; integer  $u,v$ ;
  if  $n = 0$  then
    begin  $v := 0$ ;  $u := 1$  end, which takes care of the case
     $b = 0$ 
  else
    begin  $v := 1$ ;  $u := 0$ ;
    for  $i := n-1$  step -1 until 1 do
      begin integer  $t$ ;
       $t := v$ ;  $v := u - v \times q[i]$ ;  $u := t$ 
      end  $i$ 
    end the case  $n \neq 0$ . It remains now to multiply the equality
     $1 = au_1 + bv_1$  through by  $c$ ;
    begin integer  $x_0, y_0$ ;
     $x_0 := c \times u \times \text{sign}(a)$ ;  $y_0 := c \times v \times \text{sign}(b)$ ; print ( $x_0, y_0$ );
    comment If  $x_0, y_0$  is a particular solution then  $x_0 \pm ib$ ,
     $y_0 \mp ia$ ,  $i=1,2, \dots$  gives the general solution. Therefore;
    if general solution then
      begin  $u := b$ ;  $v := a$ ;
       $A := \text{print}(x_0 + u, y_0 - v)$ ;  $\text{print}(x_0 - u, y_0 + v)$ ;
       $u := u + b$ ;  $v := v + a$ ;
      if time permits then go to  $A$ 
      end general solution and
    end solution.
  end  $u,v$ 
end Diophantus.

```

ALGORITHM 140  
MATRIX INVERSION  
P. Z. INGERMAN

University of Pennsylvania, Philadelphia, Penn.

```

procedure invert ( $a$ ) of order:( $n$ ) with tolerance:( $\text{eps}$ ) and
error exit:(oops);
value  $n, \text{eps}$ ; array  $a$ ; integer  $n$ ; real  $\text{eps}$ ; label oops;
comment This procedure inverts a matrix by using elementary
row operations. Although the method is not particularly good
for ill-conditioned matrices, the simplicity of the algorithm
and the fact that the inversion occurs in place make it useful
on occasion;
begin integer  $i$ ;
for  $i := 1$  step 1 until  $n$  do
  begin integer  $j, k$ ; real  $q$ ;
   $q := a[i,i]$ ;
  if  $\text{abs}(q) \leq \text{abs}(\text{eps})$  then go to oops;
   $a[i,i] := 1$ ;
  if  $q \neq 1$  then for  $k := 1$  step 1 until  $n$  do  $a[i,k] := a[i,k]/q$ ;
  for  $j := 1$  step 1 until  $n$  do
    if  $i \neq j$  then
      begin  $q := a[j,i]$ ;  $a[j,i] := 0$ ;
      for  $k := 1$  step 1 until  $n$  do
         $a[j,k] := a[j,k] - q \times a[i,k]$  end end end

```

ALGORITHM 141  
PATH MATRIX

P. Z. INGERMAN

University of Pennsylvania, Philadelphia, Penn.

```

procedure find path ( $a, n$ );
value  $n$ ; Boolean array  $a$ ; integer  $n$ ;
comment This procedure is merely an Algol implementation
of the method of Warshall (JACM 9(1962), 11-12). Some ad-
vantage is taken of the characteristics of the problem to in-
crease the efficiency;
begin integer  $i, j, k$ ;
for  $j := 1$  step 1 until  $n$  do
  for  $i := 1$  step 1 until  $n$  do
    if  $a[i,j] \wedge i \neq j$  then
      for  $k := 1$  step 1 until  $n$  do
         $a[i,k] := a[i,k] \vee a[j,k]$  end findpath

```

CERTIFICATION OF THE CALCULATION OF  
EASTER...[Donald Knuth, *Comm. A.C.M.*, Apr. 1962]

M. R. WILLIAMS

University of Alberta, Calgary, Alberta, Canada

The two programs, written to demonstrate ALGOL and COBOL, were translated into FORTRAN for the IBM 1620. Both programs correctly determined the month and day of the "Western Easter" for the years 1901 to 1999. No further checking was done because a more comprehensive reference list of the dates of the "Western Easter" was not available.

If the statement:

```

 $\text{epact} := \text{mod}(11 \times \text{golden number} + 20 + \text{Clavian}$ 
 $\text{correction} - \text{Gregorian correction}, 30)$ ;

```

is changed to:

```

 $\text{epact} := \text{mod}(11 \times \text{golden number} + 19 + \text{Clavian}$ 
 $\text{correction} - \text{Gregorian correction}, 30) + 1$ ;

```

it eliminates the statement:

```

if  $\text{epact} \leq 0$  then  $\text{epact} := \text{epact} + 30$ ;

```

CERTIFICATION OF ALGORITHM 84  
 SIMPSON'S INTEGRATION [P. E. Hennion, *Comm. ACM*, Apr. 62]

PETER G. BEHRENZ  
 Matematikmaskinnämnden, Stockholm, Sweden

SIM was successfully run on FACIT EDB using FACIT-ALGOL 1, which is a realization of ALGOL 60 for FACIT EDB. No changes in the program were necessary. To test SIM some polynomials were integrated.

CERTIFICATION OF ALGORITHM 94  
 COMBINATION [J. Kurtzberg, *Comm. ACM*, June 1962]  
 RONALD W. MAY

University of Alberta, Calgary, Alberta, Canada

Algorithm 94 was translated into FORTRAN for the IBM 1620 and run successfully with no corrections. The variable A, however, has not been declared.

REMARK ON ALGORITHM 99  
 EVALUATION OF JACOBI SYMBOL [S. J. Garland and A. W. Knapp, *Comm. ACM* 6, June 1962]

RONALD W. MAY  
 University of Alberta, Calgary, Alberta, Canada

One syntactical error was found in this procedure. It occurs in the second if statement following the label even. The statement

```
if q then if parity ((m↑2-1)÷8) then
  p := ¬ p;
```

might be changed as follows.

```
if q then go to CHECK;
next 1: if n = 1 then go to done;
CHECK: if parity ((m ↑ 2 - 1) ÷ 8) then
  p := ¬ p;
  go to next 1;
```

The two statements beginning with CHECK could be inserted before the label done and after the statement go to loop;

REMARK ON ALGORITHM 106  
 COMPLEX NUMBER TO A REAL POWER [Margaret L. Johnson and Ward Sangren, *Comm. ACM* 5, Jul. 1962]

GRANT W. ERWIN, JR.  
 The Boeing Co., Renton, Wash.

The comment "if W is a reciprocal integer it does not follow that the desired power (a root) will be calculated" might better read "if W is the reciprocal of an integer N, the procedure will calculate an nth root, but possibly not the particular nth root desired. E.g.  $w = \frac{1}{2}$ ,  $x = -1$ ,  $y = 0$  yields  $A = \frac{1}{2}$ ,  $B = \frac{1}{2}\sqrt{3}$  rather than the simpler  $A = -1$ ,  $B = 0$ ."

The comment should be made that it is assumed that the arctan function yields a result between  $-\pi/2$  and  $\pi/2$ .

The following four corrections should be made:

- (1) if  $x < 0 \wedge y < 0$  then begin THETA: = 3.1415927;  
 should read: ... THETA: = -3.1415927;
- (2) go to RETURN end;  
 should read: go to RETURN end;
- (3) if  $x = 0 \wedge y < 0$  ...  
 should read: if  $x = 0 \wedge y > 0$  ...

- (4) if  $x = 0 \wedge y > 0$   
 should read: if  $x = 0 \wedge y < 0$  ...

CERTIFICATION OF ALGORITHM 135  
 CROUT WITH EQUILIBRATION AND ITERATION  
 [William Marshall McKeeman,\* *Comm. ACM*, Nov. 1962]

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A BALGOL translation of the algorithm was tested for accuracy, proper termination and running time on the Burroughs 220. The exact inverse of the Hilbert segment of order 6 can be stored in the 8-decimal-digit floating word of the B220 and was used in the accuracy and termination tests. The Hilbert segment  $H_6$  is very ill-conditioned (for the spectral norm,  $\|H_6\| \cdot \|H_6^{-1}\| = 1.3 \times 10^7$ ). Hence the number of iterations required should not be taken as typical.

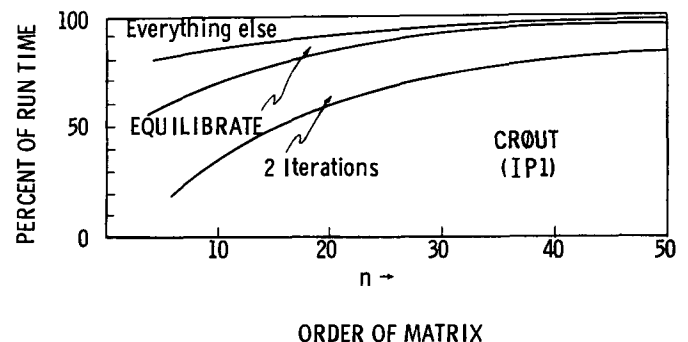
The  $[n, n]$  element (mathematically  $\frac{1}{17}$  = .090909 ...) is representative of the behavior of the rest:

	"exact" equilibration (by powers of 10)	equilibration by largest element in row
initial solution	.092587535	.094091506
first iteration	.090877240	.091498265
second iteration	.090909695	.091570311
third iteration	.090909080	.091568310
fourth iteration	.090909091	.091568365
fifth iteration	terminated	.091568364 terminated

Conclusions: The iterating procedure terminated correctly, or performed one extra iteration in each case. If the equilibration procedure alters the data, the iteration will converge to the solution for the altered matrix. If the matrix is ill-conditioned, as in the case above, the equilibration may cost a great deal more than it gains. As a practical matter, a machine language substitute for EQUILIBRATE which will not cause rounding of the data is probably the best course of action.

The running time is approximately proportional to  $n^3$  as expected. If for a given machine,  $\mu$  is the floating multiply time in seconds, one can expect that run time will be given by  $rt := 1.3 \times \mu \times (n + 7) \uparrow 3$  seconds for a call on LINEARSYSTEM with one right-hand side.

The division of run time between the various phases of the algorithm is as follows:



REFERENCE:

1. SAVAGE AND LUKACS, Tables of inverses of finite segment of the Hilbert matrix. In Olga Taussky (Ed.), Contributions to the Solution of Systems of Linear Equations and the Determination of Eigenvalues, pp. 105-108, Nat. Bur. Standards Appl. Math. Series no. 39, U. S. Government Printing Office, Wash., D.C., 1954.