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ALGORITHM 257
HAVIE INTEGRATOR [D1]
real procedure havieintegrator (x, a, b, eps, integrand, m);
 value a, b, eps, m; integer m;
 real integrand, x, a, b, eps;
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ROBERT N. Kubik (Reed. 9 June 1964 and 21 Dec. 1964) The Babcock & Wilcox Co. Lynchburg, Viriginia

comment This algorithm performs numerical integration of definite integrals using an equidistant sampling of the function and repeated halving of the sampling interval. Each halving allows the calculation of a trapezium and a tangent formula on a finer grid, but also the calculation of several higher order formulas which are defined implicitly. The two families of approximate solutions will normally bracket the value of the integral and from these convergence is tested on each of the several orders of approximation. The algorithm is based on a private communication from F. Håvie of the Institutt for Atomenergi Kjeller Research Establishment, Norway. A FORTRAN version of the algorithm is in use on the Philco-2000. A few test cases have been run on the Burroughs B5000. In particular, a and b are the lower and upper limits of integration, respectively, eps is the convergence criterion, integrand is the value of the function to be integrated (sampled), and m is the maximum order approximation to be considered in attempting to satisfy the eps convergence criterion. If convergence is not gained, then the value returned is that of the nonlocal variable, mask. The parameter integrand must be an expression involving the variable of integration x. See the driver program of this algorithm for ex-

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amples of the procedure call;
begin real h, endpts, sumt, sumu, d;
  integer i, j, k, n;
  real array t, u, tprev, uprev[1:m];
  x := a; endpts := integrand; x := b; endpts := 0.5 \times
    (integrand + endpts);
  sumt := 0.0; i := n := 1; h := b - a;
estimate: l[1] := h \times (endpts + sumt); \quad sumu := 0.0;
 comment t[1] = h \times (0.5 \times f[0] + f[1] + f[2] + \dots + 0.5 \times f[2^{i-1}]);
 x := a - h/2.0;
 for j := 1 step 1 until n do
 begin
   x := x + h; sumu := sumu + integrand
 end:
 u[1] := h \times sumu; \quad k := 1;
 comment u[1] = h \times (f[1/2] + f[3/2] + \cdots + f[(2^{i}-1)/2]), k
    corresponds to approximate solution with truncation error
    term of order 2k;
test: if abs(t[k]-u[k]) \le eps then
 begin
   havie integrator := 0.5 \times (t[k] + u[k]); go to exit
 if k \neq i then
 begin
   d := 2 \uparrow (2 \times k);
   t[k{+}1] \,:=\, (d{\times}t[k]{-}tprev[k])/(d{-}1.0)\,;
   tprev[k] := t[k];
   u[k\!+\!1] \,:=\, (d\!\times\! u[k]\!-\!uprev[k])/(d\!-\!1.0)\,;
   uprev[k] := u[k];
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comment This implicit formulation of the higher order integration formulas is given in [ROMBERG, W. Vereinfachte Numerische Integration. Det Kong. Norske Videnskabers Selskabs Forhandl. 28, 7 (1955), Trondheim; and in Stiefel, E. Einführung in der Numerische Mathematik. Teubner Verlagsges., Stuttgart, 1961, pp. 131-136. (English translation: An Introduction to Numerical Mathematics, Academic Press, New York, 1963, pp. 149-155)]. See also Algorithm 60 where the same implicit relationship is used to calculate t[k+1] only;

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k := k + 1;
     if k = m then
     begin
      havieintegrator := mask; go to exit
     end;
    go to test
  end;
  h := h/2.0; sumt := sumt + sumu;
  tprev[k] := t[k]; uprev[k] := u[k];
  i := i + 1; \quad n := 2 \times n;
  go to estimate;
exit: end havieintegrator
Following is a driver program to test havieintegrator.
begin comment First test case, y = \int_0^{\pi/2} \cos x \, dx = 1.0
     (0.999999981 as executed on the B5000), is an example of the
    higher order approximations yielding fast convergence as in
     Algorithm 60; second test case, y = \int_{0}^{4.3} e^{-x^2} dx = .8862269255
     (.8862269739 as executed on the B5000), is an example where
    this algorithm is superior to Algorithm 60 because the higher
    order approximations converge more slowly than the linear
    approximations; see also [Thacher, H. C., Jr., Remark on
    Algorithm 60. Comm. A.C.M. 7 (July 1964), 420];
  real a, b, eps, mask, y, answer;
  a := 0.0; b := 1.5707963; eps := 0.000001; mask := 9.99;
  answer := havie integrator (y, a, b, eps, cos(y), 12);
  outreal (1, answer);
  a := 0.0; b := 4.3;
  answer := havieintegrator (y, a, b, eps, exp(-y \times y), 12);
  outreal (1, answer);
end
ALGORITHM 258
TRANSPORT [H]
G. BAYER (Reed. 4 May 1964 and 4 Mar. 1965);
Technische Hochschule, Braunschweig, Germany
procedure transport (c, x, a, b, m, n, inf, cost);
  value m, n, inf; integer m, n, inf, cost;
    integer array c, x, a, b;
comment The parameters are c[i,j] array of costs, the quantities
  available a[i], the quantities required b[j], i = 1, \dots, m, j =
  1, \dots, n. Sum of a[i] = \text{sum of } b[j]. inf has to be the greatest
  positive integer within machine capacity, all quantities have to
  be integer. The flows x[i, j] are computed by the "primal-dual-
  algorithm," cited in [Hadley, G. Linear Programming. Read-
  ing, London, 1962, pp. 351-367]. The procedure follows the de-
  scription given on p. 357. Multiple solutions are left out of
  account;
begin integer i, j, p, h, k, y, t, l;
  integer array v, xsj, s, r, listv[1:n], u, xis, d, g, listu[1:m];
  Boolean array xb[1:m, 1:n];
  integer procedure sum(i, a, b, x); value a, b;
    integer i, a, b, x;
    begin integer 8;
     s := 0;
      for i := a step 1 until b do s := s + x;
     sum := s
   end;
 comment Array xb for notation of "circled cells," listu and
   lists of labeled rows and columns. Other notations follow
 for i := 1 step 1 until m do xis[i] := a[i];
 for j := 1 step 1 until n do xsj[j] := b[j];
 for i := 1 step 1 until m do
 begin h := inf; for j := 1 step 1 until n do
   begin x[i, j] := 0; p := c[i, j]; if p < h then h := p end;
   u[i] := h;
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for j := 1 step 1 until n do
        xb[i,j] := if c[i,j] = h then true else false
   end u[i];
   for j := 1 step 1 until n do
   begin h := inf;
      for i := 1 step 1 until m do
      begin if xb[i, j] then
        begin v[j] := 0; go to aa end;
        d[i] := p := c[i, j] - u[i];
        if p < h then h := p
     end;
     v[j] := h;
     for i := 1 step 1 until m do
     begin if d[i] = h then xb[i, j] := true end;
   end v[j];
   for j := 1 step 1 until n do listv[j] := 0;
   for i := 1 step 1 until m do listu[i] := 0;
 s2: for i := 1 step 1 until m do
   begin for j := 1 step 1 until n do
     begin if xb[i, j] then
       begin h := x[i, j] := \mathbf{if} \ xsj[j] \le xis[i]
          then xsj[j] else xis[i];
         xsj[j] := xsj[j] - h;
         xis[i] := xis[i] - h
       end
     end
   end:
 s03: if sum(j, 1, n, xsj[j]) = 0 then go to s6;
   for j := 1 step 1 until n do s[j] := r[j] := 0;
  h := 0; k := 1;
 s3: for i := 1 step 1 until m do
   begin if xis[i] > 0 then
     begin d[i] := xis[i]; g[i] := 2 \times n;
       for j := 1 step 1 until n do
       begin if xb[i, j] \wedge r[j] = 0 then
         begin s[j] := d[i]; r[j] := i; listv[k] := j; k := k + 1;
           if xsj[j] > h then
           \mathbf{begin}\ h\ :=\ xsj[j];\quad p\ :=\ j\ \mathbf{end}
         end
       end
    end
    else d[i] := g[i] := 0
  end:
s53: if k = 1 then go to s13;
  l := 1;
  for k := 1 step 1 until n do
  begin j := listv[k]; listv[k] := 0; if <math>j = 0 then go to $33;
    for i := 1 step 1 until m do
    begin if xb[i,j] \land x[i,j] > 0 \land g[i] = 0 then
      \mathbf{begin}\ d[i]\ :=\ \mathbf{if}\ x[i,j]\ \leq\ s[j]
         then x[i, j] else s[j];
         g[i] := j; listu[l] := i; l := l + 1
      end
    end
  end;
s33: if l = 1 then go to s13:
 k := 1;
  for l := 1 step 1 until m do
  begin i := listu[l]; listu[l] := 0; if <math>i = 0 then go to s43;
    for j := 1 step 1 until n do
   begin if xb[i, j] \wedge r[j] = 0 then
      \mathbf{begin}\ s[j]\ :=\ d[i];\ \ r[j]\ :=\ i;\ \ listv[k]\ :=\ j;\ \ k\ :=\ k\ +\ 1;
        if xsj[j] > h then
        begin h := xsj[j]; p := j end
     end
   end
 end;
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s43: go to s53;
$13:; comment end of labeling process;
  if h > 0 then go to 84 else
    if sum(j, 1, n, xsj[j]) = 0 then go to s6 else go to s5:
s4: k := p;
  h := \mathbf{if} \ s[k] < xsj[k] \ \mathbf{then} \ s[k] \ \mathbf{else} \ xsj[k]:
s41: y := r[k]; x[y, k] := x[y, k] + h;
  xis[y] := xis[y] - h; \quad xsj[k] := xsj[k] - h;
  t := g[y]; if t = 2 \times n then go to s03; x[y, t] := x[y, t] - h:
  xis[y] := xis[y] + h;  xsj[t] := xsj[t] + h;  k := t;  go to st]
s5: h := inf;
  for i := 1 step 1 until m do
  for j := 1 step 1 until n do
  begin if g[i] \neq 0 \land r[j] = 0 then
    begin p := c[i, j] - u[i] - v[j];
      \mathbf{if}\ p\ < h\ \mathbf{then}\ h\ :=\ p
    end
  end;
  for i := 1 step 1 until m do
  begin if g[i] \neq 0 then u[i] := u[i] + h end;
  for j := 1 step 1 until n do
  begin if r[j] \neq 0 then v[j] := v[j] - h end;
  for i := 1 step 1 until m do
  for j := 1 step 1 until n do
  begin if c[i, j] = u[i] + v[j] then xb[i, j] := true end;
s6: cost := sum(i, 1, m, a[i] \times u[i]) + sum(j, 1, n, b[j] \times v[j])
end;
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CERTIFICATION OF ALGORITHM 246 [Z] GRAYCODE [J. Boothroyd, Comm. ACM 7 (Dec. 1964), 701]

WILLIAM D. ALLEN (Recd. 8 Feb. 1965 and 23 Feb. 1965) Computing Ctr., U. of Kentucky, Lexington, Ky.

graycode was coded in Fortran IV and tested on the IBM 7040. graycode code was generated from 0 to 10,000 in both ascending and descending sequence. The procedure required no corrections and gave correct results for all cases tested.

ERRATUM

In the paper "Wengert's Numerical Method for Partial Derivatives, Orbit Determination and Quasilinearization," by R. E. Bellman, H. Kagiwada and R. E. Kalaba, Comm. ACM 8 (Apr. 1965), the authors submit the following erratum.

After Equation (7), the line should read:

where the minimization is over $x_{k+1}(t_1)$, $\dot{x}_{k+1}(t_1)$, $y_{k+1}(t_1)$ and $\dot{y}_{k+1}(t_1)$.