ALGORITHM 283

SIMULTANEOUS DISPLACEMENT OF POLYNO-MIAL ROOTS IF REAL AND SIMPLE [C2]

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procedure Prrs (A, X, n, eps); value n, eps; integer n; real eps; array A, X;

comment Prrs (polynomial roots real simple) computes the n roots X of the polynomial equation

$$A_n x^n + A_{n-1} x^{n-1} + \cdots + A_0 = 0$$

simultaneously. On entry the array X contains the vector of initial approximations to the roots and on exit it contains the vector of improved approximations to the roots. The initial approximations must be distinct. Accuracy is specified by means of a parameter eps. Iteration is continued until the Euclidean norm of the correction vector does not exceed eps. The convergence is quadratic;

begin integer i, k; real x, P, Q; $eps := eps \uparrow 2$; $W \colon Q := 0$;

for i := 1 step 1 until n do

begin x := P := A[n];

for k := 1 step 1 until n do

begin $x := x \times X[i] + A[n-k]$;

if $k \neq i$ then $P := P \times (X[i] - X[k])$ end; X[i] := X[i] - x/P; $Q := Q + (x/P) \uparrow 2$ end;

if Q > eps then go to Wend

CERTIFICATION OF ALGORITHM 9 [D2]

RUNGE-KUTTA INTEGRATION [P. Naur et al., Comm. ACM 3 (May 1960), 318]

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Algorithm 9 was transcribed into the hardware representation for CDC 3600 Algoriand run successfully. The following procedure was used for the global procedure comp:

real procedure $comp\ (a,b,c)$; value a,b,c; real a,b,c; begin integer $AE,\ BE,\ CE;$

integer procedure expon(x); real x;

comment This function produces the base 10 exponent of x; $expon := \mathbf{if} x = 0$ then -999 else

entier $(.4342944819 \times ln(abs(x)) + 1)$;

comment The number -999 may be replaced by any number less than the exponent of the smallest positive number handled by the particular machine used, for this algorithm assumes that true zero has an exponent smaller than any nonzero floating-point number. Users implementing real procedure comp by machine code should make sure that this condition is satisfied by their program;

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AE := expon(a); \quad BE := expon(b); \quad CE := expon(c); if AE < BE then AE := BE; if AE < CE then AE := CE; comp := abs(a-b)/10 \uparrow AE nd
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This has the advantage of machine independence, but is highly inefficient compared to machine code.

The procedure was tested using the two following procedures for FKT:

procedure FKT (X, Y, N, Z); real X; integer N; array Y, Z;

comment $(dy_1/dx) = z_1 = y_2$, $(dy_2/dx) = z_2 = -y_1$. With $y_1(0) = 0$, $y_2(0) = 1$, the solution is $y_1 = \sin x$, $y_2 = \cos x$;

begin Z[1] := Y[2]; Z[2] := -Y[1] **end**;

procedure FKT(X, Y, N, Z); real X; integer N; array Y, Z;

comment $(dy_1/dx) = 1 + y_1^2$. For $y_1(0) = 0$, $y(x) = \tan x$; $Z[1] := 1 + Y[1] \uparrow 2$;

The RK procedure was used to integrate the differential equations represented by the first FKT procedure from x=0(0.5)7.0, with $eps=eta=10^{-6}$, and with $y_1(0)=0$, $y_2(0)=1$. The actual step size h was .0625 for most of the range, but was reduced to .03125 in the neighborhood of $x=k\pi/2$, where one or the other of the solutions is small.

The computed solutions at x=7.0 were: $y_1=6.5698602746 \times 10^{-1}$, $y_2=7.5390270246 \times 10^{-1}$, with errors -5.71×10^{-7} and 4.48×10^{-7} , respectively.

Results for the second differential equation are summarized in Table I below.

The efficiency of the procedure would be increased slightly on most computers by changing the type of the **own** variable s from real to integer.

The error is estimated by comparing the results of successive pairs of steps with that of a single double step. This is somewhat more time-consuming than the Kutta-Merson process presented in Algorithm 218 [Comm. ACM 6 (Dec. 1963) 737-8]. However, the criterion for step-size variation in Algorithm 9 which effectively applies an approximate relative error criterion, eps, for |y| > eta, and an absolute error criterion $eta \times eps$, for |y| < eta, appears superior when the solution fluctuates in magnitude.

REMARK ON ALGORITHM 218 [D2]

KUTTA-MERSON [Phyllis M. Lukehart, Comm. ACM 6 (Dec. 1963), 737]

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Successive calls of *Kutta Merson* with $first \equiv false$ do not reach the upper bound t+h if the interval h is unequal to the interval h of the first call with $first \equiv true$.

Proposed correction:

- 1) declaration real hc, instead of own real hc;
- 2) if first then begin for i := 1 step 1 until n do y0[i] := y[i]; hc := h; ploc := 1; first := false end else hc := h/ploc;

instead of if first then begin · · · end;

TABLE I [ALG. 9]

	η	x = 0.5			x = 1.0			x = 1.5		
		h_{min}	Absolute error	Relative error	hmin	Absolute error	Relative error	h_{min}	Absolute error	Relative error
$ \begin{array}{r} 10^{-7} \\ 10^{-5} \\ 10^{-3} \end{array} $	$ \begin{array}{c c} 10^{-3} \\ 10^{-3} \\ 10^{-3} \end{array} $.03125 .125 .25	$ \begin{array}{ c c c c c } -1 \times 10^{-9} \\ -5 \times 10^{-7} \\ -1 \times 10^{-5} \end{array} $	-9×10^{-7}	.03125 .0625 .25	$ \begin{array}{ c c c c c c } \hline 9 \times 10^{-8} \\ 8 \times 10^{-7} \\ -2 \times 10^{-4} \end{array} $	$ \begin{array}{c} 6 \times 10^{-8} \\ 5 \times 10^{-7} \\ -1 \times 10^{-4} \end{array} $.0078125	$ \begin{array}{ c c c c c } \hline -1 \times 10^{-6} \\ -2 \times 10^{-4} \\ -3 \times 10^{-2} \end{array} $	$ \begin{array}{c c} -8 \times 10^{-8} \\ -1 \times 10^{-6} \\ -2 \times 10^{-3} \end{array} $