

Algorithms

J. G. HERRIOT, Editor

ALGORITHM 352 CHARACTERISTIC VALUES AND ASSOCIATED SOLUTIONS OF MATHIEU'S DIFFERENTIAL EQUATION [S22]

DONALD S. CLEMM (Recd. 2 June
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Aerospace Research Laboratories
Wright-Patterson Air Force Base
OH 45433

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Comments Algorithm 352 is a package of
double-precision FORTRAN routines which
consists of the following primary routines:

MFCVAL—referred to as Algorithm 352
(Part A)

MATH—referred to as Algorithm 352
(Part B)

BESSEL—referred to as Algorithm 352
(Part C)

MFCVAL computes characteristic values of
Mathieu's differential equation. MATH
computes the associated solutions of this
equation, using BESSEL as an auxiliary
routine to evaluate Bessel functions. This
latter routine may be used independently.

There are other, secondary routines in-
cluded in the package, and the numbering
system (e.g. Algorithm 352 (Part A.1)) in-
dicates somewhat the mutual relation between
them, as well as their relation to the primary
routines. The functioning of the routines
and the linkages between them are explained
in the comments prefacing each one. All
literature citations refer to the following
list.

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Algorithm 352 (Part A)
MFCVAL (Characteristic Values)

Comments The subroutine MFCVAL com-
putes the first N characteristic values, a , to-
gether with upper and lower bounds, of
Mathieu's differential equation for nonnega-
tive values of the real parameter, q . The
equation can be written in the form

$$y'' + (a - 2q \cos 2x)y = 0, \quad (1)$$

where $a = a_r$ ($a = b_r$) indicates a character-
istic value associated with the even (odd)
periodic solutions.

The method consists of three steps: (1)
calculate a rough approximation based on
coefficients obtained from curve-fitting of
available tabulations, (2) determine crude
upper and lower bounds, and (3) iterate,
using a variation of Newton's method. For
a justification of this method, see [3].

Explanation of the arguments:

- N the given number of characteristic
values desired
R given as N-1 or N according as the
characteristic values are to be asso-
ciated with the even or odd solu-
tions, respectively
QQ the given nonnegative parameter q
CV the computed 6 by N array of charac-
teristic values and bounds
J the number of characteristic values
successfully computed. $J \neq N$ indi-
cates that J values were computed

with an error occurring on the $J + 1$
value. A printed message will ac-
company such an error condition.

The output array, CV, must be appro-
priately dimensioned in the calling program
and upon return will contain the following
data:

- For the Kth characteristic value, $K = 1,$
 $2, \dots, J,$
CV (1, K) the characteristic value a
CV (2, K) the function $D(a) = -T_m(a)/$
 $T_m'(a)$
CV (3, K) a_L , a lower bound of a
CV (4, K) the function $D(a_L)$
CV (5, K) a_U , an upper bound of a
CV (6, K) the function $D(a_U)$.

Reference is again made to [3], where the
function $T_m(a)$ is defined and it is proved
that $T_m(a) = 0$ if and only if a is a charac-
teristic value. From this, it can be said that
the function D is an indication of the ac-
curacy of its argument, since $a + D(a)$
would be the value of the next iteration.

The first executable statement in
MFCVAL sets a tolerance of 10^{-12} . This may
be changed by the user, but the following
comments should be heeded if it is at-
tempted.

If it is desired to reduce the tolerance in
order to achieve the greatest possible ac-
curacy, care should be taken that the toler-
ance is not less than $10^{-(n-2)}$ when execut-
ing the routines on a machine which uses n -digit
arithmetic. In other words, if the user's
computer employs 24-digit arithmetic, this
tolerance should be no less than 10^{-22} . A too
small tolerance will impose an unattainable
accuracy requirement and overflow may
occur.¹

On the other hand, some time-saving may
be achieved, at the expense of accuracy, by
making the tolerance less stringent. A toler-
ance of 10^{-d} will produce results good to at
least d digits. This is a conservative esti-
mate, since one additional iteration is per-
formed after the tolerance is met and, nor-
mally, the convergence of successive itera-
tions is quadratic.

Perhaps it should be noted again that the
accuracy of any characteristic value, a , can
be determined from the size of it relative to
the function $D(a)$. See the description of the

¹ The constant in statement numbers 425 and 445 is intro-
duced to avoid the possibility of a zero tolerance. This
should not be altered unless the routines are being run on
a machine which uses arithmetic of more than 16 digits,
and then it must not be less than $10^{-(n-2)}$, with n defined
as above.

contents of the output array CV. MFCVAL calls on the subroutines:

BOUNDS—referred to as Algorithm 352 (Part A.1)

MFITR8—referred to as Algorithm 352 (Part A.2)

TMOFA—referred to as Algorithm 352 (Part A.3)

```

C SUBROUTINE MFCVAL (N,R,QQ,CV,J)
*****
INTEGER
* J,K,KK,L,M,N,R,TYPE
DOUBLE PRECISION
* A,CV,DL,DR,DTM,Q,QQ,
* T,TM,TOL,TOLA
DIMENSION
* CV(6,N)
EQUIVALENCE
* (DL,DR,T)
COMMON /MF1/
* Q,TOL,TYPE,DUMMY(4)
TOL = 1.D-13
IF (N-R) 10,10,20
10 L = 1
20 L = 2
30 Q = QQ
DO 500 K = 1,N
J = K
IF (Q) 960,490,40
40 KK = MIN0(K,4)
TYPE = 2*MOD(L,2)+MOD(K-L+1,2)
C FIRST APPROXIMATION
GO TO (100,200,300,400), KK
100 IF (Q-1.D0) 110,140,140
110 GO TO (120,130), L
120 A = 1.D0-Q-.125D0*Q*Q
130 A = Q*Q
A = A*(-.5D0+.0546875D0*A)
140 IF (Q-2.D0) 150,180,180
150 GO TO (160,170), L
160 A = 1.033D0-1.0746D0*Q-
*.0688D0*Q*Q
170 A = .23D0-.495D0*Q-
*.191D0*Q*Q
180 A = -.25D0-2.D0*Q+
* 2.D0*DSQRT(Q)
200 DL = L
IF (Q*DL-6.D0) 210,350,350
210 GO TO (220,230), L
220 A = 4.01521D0-Q*
*.046D0+.0667857D0*Q)
230 A = 1.D0+1.05007D0*Q-
*.180143D0*Q*Q
300 IF (Q-8.D0) 310,350,350
310 GO TO (320,330), L
320 A = 8.93867D0+.178156D0*Q-
*.0252132D0*Q*Q
330 A = 3.70017D0+.953485D0*Q-
*.0475065D0*(Q
350 DR = K-1
A = CV(1,K-1)-DR+

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```

* 4.D0*DSQRT(Q)
GO TO 420
400 A = CV(1,K-1)-CV(1,K-2)
A = 3.D0*A+CV(1,K-3)
420 IF (Q.GE.1.D0) GO TO 440
IF (K.NE.1) GO TO 430
425 TOLA = DMAX1(DMIN1(TOL,DABS(A))
* .1.D-14)
430 TOLA = TOL*DABS(A)
440 TOLA = TOL*DMAX1(Q,DABS(A))
445 TOLA = DMAX1(DMIN1(TOLA,DABS(A))
* .4D0*DSQRT(Q))
* .1.D-14)
C CRUDE UPPER AND LOWER BOUNDS
450 CALL BOUNDS (K,A,TOLA,CV,N,M)
IF (M.NE.0)
IF (M-1) 470,910,900
C ITERATE
CALL MFITR8 (TOLA,CV(1,K),
* CV(2,K),M)
IF (M.GT.0) GO TO 920
C FINAL BOUNDS AND FUNCTIONS, D
470 T = CV(1,K)-TOLA
CALL TMOFA (T,TM,DTM,M)
IF (M.GT.0) GO TO 940
CV(3,K) = T
CV(4,K) = -TM/DTM
480 T = CV(1,K)+TOLA
CALL TMOFA (T,TM,DTM,M)
IF (M.GT.0) GO TO 950
CV(5,K) = T
CV(6,K) = -TM/DTM
GO TO 500
C Q EQUALS ZERO
490 CV(1,K) = (K-L+1)**2
CV(2,K) = 0.D0
CV(3,K) = CV(1,K)
CV(4,K) = 0.D0
CV(5,K) = CV(1,K)
CV(6,K) = 0.D0
500 CONTINUE
550 RETURN
C PRINT ERROR MESSAGES
900 WRITE (6,901) K
901 FORMAT(20HO CRUDE BOUNDS CANNOT,
* 22H BE LOCATED, NO OUTPUT,
* 7H FOR K=12)
GO TO 930
910 WRITE (6,911) K
911 FORMAT(20HO ERROR IN SUBPROGRAM,
* 22H TMOFA, VIA SUBPROGRAM,
* 18H BOUNDS, NO OUTPUT,
* 7H FOR K=12)
GO TO 930
920 WRITE (6,921) K
921 FORMAT(20HO ERROR IN SUBPROGRAM,
* 22H TMOFA, VIA SUBPROGRAM,
* 18H MFITR8, NO OUTPUT,
* 7H FOR K=12)
930 J = J-1
GO TO 550
940 WRITE (6,941) K
941 FORMAT(20HO ERROR IN SUBPROGRAM,
* 22H TMOFA, NO LOWER BOUND,
* 7H FOR K=12)
CV(3,K) = 0.D0
CV(4,K) = 0.D0
GO TO 480
950 WRITE (6,951) K
951 FORMAT(20HO ERROR IN SUBPROGRAM,
* 22H TMOFA, NO UPPER BOUND,
* 7H FOR K=12)
CV(5,K) = 0.D0
CV(6,K) = 0.D0
GO TO 500
960 WRITE (6,961)
961 FORMAT(20HO Q GIVEN NEGATIVELY,,
* 20H USED ABSOLUTE VALUE)
Q = -Q
GO TO 40
END

```

Algorithm 352 (Part B)
MATH (Mathieu Functions)

Comments The subroutine MATH computes various solutions (and their derivatives), of either Mathieu's differential equation or Mathieu's modified equation, which are associated with the characteristic values.

The even periodic solution of equation (1) is

$$ce_r(x,q) = \sum_{k=0}^{\infty} A_{2k+p} \cos(2k+p)x, \quad (2)$$

associated with $a_r(q)$, and the odd periodic solution is

$$se_r(x,q) = \sum_{k=0}^{\infty} B_{2k+p} \sin(2k+p)x, \quad (3)$$

associated with $b_r(q)$. The order, r , is of the form $2n + p$. The n is a nonnegative integer while $p = 0$ or 1 indicates the solution is of period π or 2π . Calculation of the periodic solutions allows the following three options of normalization:

(a) Neutral. We define neutral coefficients such that $\bar{A}_{2k+p} = A_{2k+p}/A_{2s+p}$, where s is chosen so that A_{2s+p} is the numerically largest one of the set. The \bar{B}_{2k+p} are similarly defined. This has the computationally convenient effect of making the largest coefficient equal to unity, hence all calculations are carried out with them. If a normalization other than neutral is selected, it is effected on the output array F only, the coefficients themselves remaining unchanged.

(b) Ince. The normalization adopted in [6] is defined so that if $y(x,q)$ represents either function (2) or (3) then

$$\int_0^{2\pi} y^2(x,q) dx = \pi.$$

(c) Stratton. As defined in [8], and in the notation of [7], this normalization is effected so that

$$Se_r(q,0) = \left[\frac{d}{dx} So_r(q,x) \right]_{x=0} = 1,$$

where Se is the even solution and So the odd. If we replace x by ix in (1), we get

$$y'' - (a-2q \cosh 2x) y = 0, \quad (4)$$

known as Mathieu's modified equation. The solutions of (4) have been termed radial in [8] and, for characteristic values, can be put in the following form, using the notation of [4] and [5]:

$$Mc_r^{(j)}(x,q) = \sum_{k=0}^{\infty} (-1)^{n+k} A_{2k+p} [F_k + G_k] / A_{2s+p} \epsilon_{2s+p}, \quad (5)$$

associated with $a_r(q)$, and

$$Ms_r^{(j)}(x,q) = \sum_{k=0}^{\infty} (-1)^{n+k} B_{2k+p} [F_k - G_k] / B_{2s+p}, \quad (6)$$

associated with $b_r(q)$. The order r equals $2n + p$, as in (2) and (3), and $\epsilon_m = 1$ if $m \neq 0$, but $\epsilon_0 = 2$. The choice of s is arbitrary here, but for numerical purposes we choose it in the manner described previously for *neutral* normalization. The coefficients are the same as defined in (2) and (3), while F_k and G_k involve the Bessel functions as follows:

$$F_k = J_{k-s}(u_1) Z_{k+p+s}^{(i)}(u_2), \quad (7)$$

$$G_k = J_{k+p+s}(u_1) Z_{k-s}^{(i)}(u_2), \quad (8)$$

$$u_1 = q^{\frac{1}{2}} e^{-x}, \quad u_2 = q^{\frac{1}{2}} e^x,$$

$$Z_m^{(1)}(u) = J_m(u), \quad Z_m^{(2)}(u) = Y_m(u).$$

The solutions (5)–(6) are said to be of the first or second kind depending on whether $j = 1$ or 2 in (5)–(8).

Explanation of the arguments:

- XX the given independent variable x
- QQ the given positive parameter q
- R the given order r
- CV the given characteristic value, $a_r(q)$ or $b_r(q)$
- SOL given as 1, 2, or 3 according as the desired solution is (1) radial of the first kind, (2) radial of the second kind, or (3) periodic
- FNC given as 1, 2, 3, or 4 according as the desired solution is (1) associated with b_r , (2) associated with a_r , (3) the derivative of solution (1), or (4) the derivative of solution (2)
- NORM given as 1, 2, or 3 according as the desired normalization is (1) defined as *neutral*, (2) defined by Ince, or (3) defined by Stratton. (This argument is decoded only if SOL = 3.)
- F the computed three-element array, containing: (1) the solution value, (2) the series term of largest magnitude, and (3) the last term included in the summation
- K the computed two-element array, containing: (1) the index, k , of the term in F(2), and (2) the index of the term in F(3)
- M the error indicator cell: M = 0 indicates successful execution of subprogram, M = 1 signifies an error condition explained by an accompanying printed message.

The accuracy of results (within limits) and the speed of convergence may be altered by the user. See SUM (Algorithm 352 (Part B.2)) for details.

MATH calls on the subroutines:

- COEF—referred to as Algorithm 352 (Part B.1)
- SUM—referred to as Algorithm 352 (Part B.2)
- BESSEL—referred to as Algorithm 352 (Part C)

```

CUBROUTINE MATH (XX,QQ,P,CV,SOL,
* *****
*
* INTEGER
* FNC,I,K(2),KLAST,KMAX,L,
* LL,M,MF,ML,MM,MO,M1,M2S,
* N,NORM,P,R,S,SOL,TYPE
*
* DOUBLE PRECISION
* A,AB,CV,DLAST,DUMAX,F(3),G,
* J,G,QQ,T,TOL,U1,U2,X,XX,Y
*
* EXTERNAL
* DC,DDC,DDS,DS,DPC,DPS,
* PC,PS
*
* COMMON
* J(250),Y(250),U1,U2,N,P,S,
* L,X,T,I,LL,G,DMAX,DLAST,
* KMAX,KLAST,DUM1(578),A,
* DUM2(6),MM,ML,AB(200)
*
* COMMON /MF1/
* Q,TOL,TYPE,M1,M0,M2S,MF
*
M = 0
IF (SOL.LT.1 .OR.
* SOL.GT.3 .OR.
* FNC.LT.1 .OR.
* FNC.GT.4) GO TO 40C
A = CV
Q = QQ
TOL = 1.D-13
TYPE = 2*MOD(FNC,2)+MOD(R,2)
CALL COEF (M)
IF (M) 410,10,420
10 N = R/2
P = MOD(R,2)
S = MM/2
L = ML/2
X = XX
T = 1.D0
IF (SOL.EQ.3)
* GO TO (150,160,170,180), FNC
U1 = DSQRT(Q)*DEXP(-X)
U2 = Q/U1
LL = L+S+P
C COMPUTE BESSEL FUNCTIONS
CALL BESSEL (1,U1,J,LL)
CALL BESSEL (SOL,U2,Y,LL)
C EVALUATE SELECTED FUNCTION
GO TO (50,60,70,80), FNC
50 CALL SUM (DS) GO TO 300
60 CALL SUM (DC) GO TO 300
70 CALL SUM (DDS) GO TO 300
80 CALL SUM (DDC) GO TO 300
150 CALL SUM (PS) GO TO 200
160 CALL SUM (PC) GO TO 200
170 CALL SUM (DPS) GO TO 200
180 CALL SUM (DPC) GO TO 200
200 IF (NORM-2) 300,210,250
C INCE NORMALIZATION
210 T = AB(1)**2
IF (TYPE.EQ.0) T = T+T
DO 220 I = 1,L
T = T+AB(I+1)**2
220 CONTINUE
T = DSQRT(T)
I = M0/2
IF (AB(I).LT.0.D0) T = -T
GO TO 300
C STRATTON NORMALIZATION
250 IF (TYPE.GT.1) GO TO 270

```

```

T = AB(1)
DO 260 I = 1,L
T = T+AB(I+1)
260 CONTINUE GO TO 300
270 T = DBLE(FLOAT(P))*AB(1)
DO 280 I = 1,L
T = T+AB(I+1)*
DBLE(FLOAT(2*I+P))
280 CONTINUE
300 F(1) = G/T
F(2) = DMAX/T
F(3) = DLAST/T
K(1) = KMAX
K(2) = KLAST
350 RETURN
C PRINT ERROR MESSAGES
400 WRITE (6,401)
401 FORMAT(18H0SOL OR FNC OUT OF,
* 17H RANGE, NO OUTPUT)
GO TO 450
410 WRITE (6,411)
411 FORMAT(15HMORE THAN 200 ,
* 22HCOEFFICIENTS REQUIRED,,
* 20H QQ AND R TOO LARGE,,
* 10H NO OUTPUT)
GO TO 450
420 WRITE (6,421)
421 FORMAT(20H0ERROR IN SUBPROGRAM,
* 22H TMOFA, VIA SUBPROGRAM,,
* 13H COEF, VERIFY,
* 21H ARGUMENTS, NO OUTPUT)
450 M = 1
F(1) = 0.D0
F(2) = 0.D0
F(3) = 0.D0
K(1) = 0
K(2) = 0
GO TO 350
END

```

Algorithm 352 (Part A.1)
BOUNDS (Crude Bounds)
(Called by MFCVAL)

Comments The subroutine BOUNDS determines crude upper and lower bounds for the K th characteristic value, $K \leq N$.

Explanation of the other arguments:

- APPROX the first approximation
- TOLA the tolerance determined by subroutine MFCVAL
- CV the 6 by N array described in subroutine MFCVAL
- N variable dimension of the CV array
- MM an indicator cell used to communicate unusual and error conditions to subroutine MFCVAL

The output, $a_0 < a < a_1$, is put into the common block labeled MF2.

BOUNDS calls on the subroutine: TMOFA—referred to as Algorithm 352 (Part A.3)

```

SUBROUTINE BOUNDS (K,APPROX,
* TOLA,CV,N,MM)
C *****
*
* INTEGER
* K,KA,M,MM,N
*
* DOUBLE PRECISION
* A,APPROX,A0,A1,CV,DTM,
* D0,D1,Q,TM,TOLA
*
* DIMENSION
* CV(6,N)
*
* COMMON /MF1/
* Q,DUMMY(7)

```

```

COMMON /MF2/
*   AO,A,A1

KA = 0

IF (K.EQ.1) GO TO 20

IF (APPROX-CV(1,K-1)) 10,10,20

10 AO = CV(1,K-1)+1.00           GO TO 30
20 AO = APPROX
30 CALL TMOFA (AO, TM, DTM, M)

IF (M.GT.0) GO TO 250

DO = -TM/DTM

IF (DO) 100,300,50

C   AO IS LOWER BOUND,
C   SEARCH FOR UPPER BOUND
50 A1 = AO+DO+.100
   CALL TMOFA (A1, TM, DTM, M)

IF (M.GT.0) GO TO 250

D1 = -TM/DTM

IF (D1) 200,350,60

60 AO = A1
   DO = D1
   KA = KA+1

IF (KA-4) 50,400,400

C   A1 IS UPPER BOUND,
C   SEARCH FOR LOWER BOUND
100 A1 = AO
    D1 = DO
    AO = DMAX1(A1+D1-.100,-2.00*Q)

IF (K.EQ.1) GO TO 110

IF (AO-CV(1,K-1)) 150,150,110

110 CALL TMOFA (AO, TM, DTM, M)

IF (M.GT.0) GO TO 250

DO = -TM/DTM

IF (DO) 120,300,200

120 KA = KA+1

IF (KA-4) 100,400,400

150 KA = KA+1

IF (KA-4) 160,400,400

160 AO = A1+DMAX1(TOLA,DABS(D1))
    GO TO 30
200 A = .500*(AO+DO+A1+D1)

IF (A.LE.A0 .OR.
*   A.GE.A1) A = .500*(AO+A1)

250 MM = M           RETURN

300 CV(1,K) = AO
310 CV(2,K) = 0.00
    M = -1

350 CV(1,K) = A1           GO TO 250
400 M = 2           GO TO 310

END           GO TO 250

```

Algorithm 352 (Part A.2)

MFITR8 (Improves Characteristic Value)
(Called by MFCVAL)

Comments Given $a_0 < a < a_1$, where a_0 is a lower and a_1 an upper bound, the subroutine MFITR8 iterates to the characteristic value, replacing one of the bounds with a better approximation at each step. The

process terminates after 40 iterations unless one of the following conditions occurs first:
(1) $a - a_0 \leq \text{TOLA}$, (2) $a_1 - a \leq \text{TOLA}$, or
(3) $|D(a)| < \text{TOLA}$. See Appendix 3, method 2, of [3] for a detailed description of this process.

Explanation of output:
CV the characteristic value, a
DCV the function $D(a)$
MM an indicator cell used to communicate an error condition to subroutine MFCVAL.
MFITR8 calls on the subroutine:
TMOFA—referred to as Algorithm 352 (Part A.3)

```

SUBROUTINE MFITR8 (TOLA, CV, DCV, MM)
*****
INTEGER
*   M, MM, N

DOUBLE PRECISION
*   A, AO, A1, A2, CV, D, DCV, DTM,
*   TM, TOLA

LOGICAL
*   LAST

COMMON /MF2/
*   AO, A, A1

N = 0
LAST = .FALSE.
50 N = N+1
   CALL TMOFA (A, TM, DTM, M)

IF (M.GT.0) GO TO 400

D = -TM/DTM

C   IS TOLERANCE MET
IF (N .EQ. 40 .OR.
*   A-A0 .LE. TOLA .OR.
*   A1-A .LE. TOLA .OR.
*   DABS(D) .LT. TOLA) LAST = .TRUE.

IF (D) 110,100,120

100 CV = A
    DCV = 0.00           GO TO 320

C   REPLACE UPPER BOUND BY A
110 A1 = A           GO TO 200

C   REPLACE LOWER BOUND BY A
120 AO = A
    A2 = A+D

IF (LAST) GO TO 300
IF (A2.GT.A0.AND.A2.LT.A1)
*   GO TO 250

A = .500*(AO+A1)           GO TO 50

250 A = A2           GO TO 50

300 IF (A2.LE.A0.OR.A2.GE.A1)
*   GO TO 350

CALL TMOFA (A2, TM, DTM, M)

IF (M.GT.0) GO TO 400

D = -TM/DTM
CV = A2
310 DCV = D
320 MM = M           RETURN

350 CV = A           GO TO 310

400 CV = 0.00
    DCV = 0.00           GO TO 320

END

```

Algorithm 352 (Part A.3)
TMOFA (Accuracy Indicator)
(Called by MFCVAL, BOUNDS, MFITR8 and COEF)

Comments The subroutine TMOFA evaluates the function $T_m(a)$ and its derivative $dT_m(a)/da$. See [3] for the definitions, theorems, and numerical methods relating to the computation of these quantities.

Explanation of the arguments:
ALFA the given argument, a
TM $T_m(a)$
DTM $dT_m(a)/da$
ND internal error indicator cell
TMOFA calls no other subprograms.

```

SUBROUTINE TMOFA (ALFA, TM, DTM, ND)
*****
INTEGER
*   K, KK, KT, L, MF, MO, M1, M2S,
*   ND, TYPE

DOUBLE PRECISION
*   A, AA, ALFA, B, DG, DTM, DTYPE,
*   F, FL, G, H(200), HP, Q, QINV,
*   Q1, Q2, T, TM, TOL, TT, V

COMMON
*   G(200,2), DG(200,2), AA,
*   A(3), B(3), DTYPE, QINV, Q1,
*   Q2, T, TT, K, L, KK, KT

COMMON /MF1/
*   Q, TOL, TYPE, M1, MO, M2S, MF

EQUIVALENCE
*   (H(1), G(1,1)), (Q1, HP),
*   (Q2, F)

DATA FL /1.0D+30/

C   STATEMENT FUNCTION
V(K) = (AA-DBLE(FLOAT(K))**2)/Q

ND = 0
KT = 0
AA = ALFA
DTYPE = TYPE
QINV = 1.00/Q
DO 10 L = 1, 2
DO 5 K = 1, 200
G(K,L) = 0.00
DG(K,L) = 0.00
5 CONTINUE
10 CONTINUE

IF (MOD(TYPE,2)) 20,30,20

20 MO = 3           GO TO 40

30 MO = TYPE+2
40 K = .500*DSQRT(DMAX1(
*   3.00*Q+AA, 0.00))
M2S = MINO(2*K+MO+4,
*   398+MOD(MO,2))

C   EVALUATION OF THE TAIL OF A
C   CONTINUED FRACTION
A(1) = 1.00
A(2) = V(M2S+2)
B(1) = V(M2S)
B(2) = A(2)*B(1)-1.00
Q1 = A(2)/B(2)
DO 50 K = 1, 200
MF = M2S+2+2*K
T = V(MF)
A(3) = T*A(2)-A(1)
B(3) = T*B(2)-B(1)
Q2 = A(3)/B(3)

*   IF (DABS(Q1-Q2) .LT. TOL)
*       GO TO 70

Q1 = Q2
A(1) = A(2)
A(2) = A(3)
B(1) = B(2)
B(2) = B(3)

```

```

50 CONTINUE
KT = 1
70 T = 1.00/T
TT = -T*T*QINV
L = MF-M2S
DO 80 K = 2,L,2
T = 1.00/(V(MF-K)-T)
TT = T*T*(TT-QINV)
80 CONTINUE
KK = M2S/2+1

IF (KT.EQ.1) Q2 = T
G(KK,2) = .5D0*(Q2+T)
DG(KK,2) = TT

C STAGE 1
G(2,1) = 1.00
DO 140 K = M0,M2S,2
KK = K/2+1

* IF (K.LT.5)
  IF (K-3) 100,110,120

G(KK,1) = V(K-2)-1.00/G(KK-1,1)
DG(KK,1) = QINV+DG(KK-1,1)/
  G(KK-1,1)**2
* GO TO 130
100 G(2,1) = V(0)
DG(2,1) = QINV GO TO 130
110 G(2,1) = V(1)+DTYPE-2.00
DG(2,1) = QINV GO TO 130
120 G(3,1) = V(2)+(DTYPE-2.00)/
  G(2,1)
* DG(3,1) = QINV+(2.00-DTYPE)*
  DG(2,1)/G(2,1)**2
* IF (TYPE.EQ.2) G(2,1) = 0.00
130 IF (DABS(G(KK,1)).LT.1.00)
  GO TO 200
140 CONTINUE

C BACKTRACK
TM = G(KK,2)-G(KK,1)
DTM = DG(KK,2)-DG(KK,1)
M1 = M2S
KT = M2S-M0
DO 180 L = 2,KT,2
K = M2S-L
KK = K/2+1
G(KK,2) = 1.00/(V(K)-G(KK+1,2))
DG(KK,2) = -G(KK,2)**2*
  (QINV-DG(KK+1,2))

* IF (K-2) 150,150,160
150 G(2,2) = 2.00*G(2,2)
DG(2,2) = 2.00*DG(2,2)
160 T = G(KK,2)-G(KK,1)
* IF (DABS(T)-DABS(TM))
  170,180,180
170 TM = T
DTM = DG(KK,2)-DG(KK,1)
M1 = K
180 CONTINUE GO TO 320

C STAGE 2
200 M1 = K
K = M2S
KK = K/2+1
210 IF (K.EQ.M1)
* IF (K-2) 300,300,310

K = K-2
KK = KK-1
T = V(K)-G(KK+1,2)
IF (DABS(T)-1.00) 250,220,220
220 G(KK,2) = 1.00/T
DG(KK,2) = (DG(KK+1,2)-QINV)/T**2
GO TO 210

C STAGE 3
250 IF (K.EQ.M1) IF (T) 220,290,220
HP = DG(KK+1,2)-QINV
260 G(KK,2) = FL
H(KK) = T

```

```

K = K-2
KK = KK-1
F = V(K)*T-1.00
IF (K.EQ.M1) IF (F) 280,290,280
IF (DABS(F)-DABS(T)) 270,280,280
270 HP = HP/T**2-QINV
T = F/T GO TO 260
280 G(KK,2) = T/F
DG(KK,2) = (HP-QINV*T*T)/F**2
GO TO 210
290 ND = 1 GO TO 320

C CHAINING M EQUALS 2
300 G(2,2) = 2.00*G(2,2)
DG(2,2) = 2.00*DG(2,2)
310 TM = G(KK,2)-G(KK,1)
DTM = DG(KK,2)-DG(KK,1)
320 RETURN
END

```

Algorithm 352 (Part B.1)
COEF (Coefficients)
(Called by MATH)

Comments The subroutine COEF computes the *neutral* coefficients, as defined in the *Comments* of Algorithm 352 (Part B), and returns them via common array AB. Argument M is an internal error indicator cell. For details of the method used, see Appendix 6 of [3]. COEF calls on the subroutine: TMOFA—referred to as Algorithm 352 (Part A.3)

```

SUBROUTINE COEF (M)
*****
INTEGER
* K,KA,KB,KK,M,MF,ML,MM,
* M0,M1,M2S,TYPE

DOUBLE PRECISION
* A,AB,FL,G,H(200),Q,T,
* TOL,V,V2

COMMON
* G(200,2),DUM1(800),A,T,K,
* KA,KB,KK,MM,ML,AB(200)

COMMON /MF/
* Q,TOL,TYPE,M1,M0,M2S,MF

EQUIVALENCE
* (H(1),G(1,1))

DATA FL,V2/1.0D+30,1.0D-15/

C STATEMENT FUNCTION
V(K) = (A-DBLE(FLOAT(K))**2)/Q

CALL TMOFA (A,T,T,M)

IF (M.NE.0) GO TO 300

DO 60 K = 1,200
AB(K) = 0.00
60 CONTINUE
KA = M1-M0+2
DO 90 K = 2,KA,2
KK = (M1-K)/2+1

IF (K-2) 70,70,80
70 AB(KK) = 1.00 GO TO 90
80 AB(KK) = AB(KK+1)/G(KK+1,1)
90 CONTINUE
KA = 0
DO 130 K = M1,M2S,2
KK = K/2+1
ML = K

```

```

IF (G(KK,2).EQ.FL) GO TO 100
AB(KK) = AB(KK-1)*G(KK,2)
GO TO 110
100 T = AB(KK-2)
IF (K.EQ.4.AND.M1.EQ.2) T = T+T
AB(KK) = T/(V(K-2)*H(KK)-1.00)
110 IF (DABS(AB(KK)).GE.1.0D-17)
  * KA = 0
IF (KA.EQ.5) GO TO 260
KA = KA+1
130 CONTINUE
T = DLOG(DABS(AB(KK))/V2)/
* DLOG(1.00/DABS(G(KK,2)))
KA = 2*IDINT(T)
ML = KA+2*M2S
IF (ML.GT.399) GO TO 400
KB = KA+2*MF
T = 1.00/V(KB)
KK = MF-M2S
DO 150 K = 2,KB,2
T = 1.00/(V(KB-K)-T)
150 CONTINUE
KK = ML/2+1
G(KK,2) = T
DO 200 K = 2,KA,2
KK = (ML-K)/2+1
G(KK,2) = 1.00/(V(ML-K)-
  G(KK+1,2))
* 200 CONTINUE
KA = M2S+2
DO 250 K = KA,ML,2
KK = K/2+1
AB(KK) = AB(KK-1)*G(KK,2)
250 CONTINUE

C NEUTRAL NORMALIZATION
260 T = AB(1)
MM = MOD(TYPE,2)
KA = MM+2
DO 280 K = KA,ML,2
KK = K/2+1
* IF (DABS(T)-DABS(AB(KK)))
  270,280,280
270 T = AB(KK)
MM = K
280 CONTINUE
DO 290 K = 1,KK
AB(K) = AB(K)/T
290 CONTINUE
300 RETURN
400 M = -1 GO TO 300
END

```

Algorithm 352 (Part B.2)
SUM (Series Evaluation)
(Called by MATH)

Comments The subroutine SUM performs the summation, truncating the series when the magnitude of two successive terms, relative to the magnitude of the largest term, is less than or equal to 10^{-13} .

If the user is willing to accept reduced accuracy, he may save some computing time by making this tolerance larger. On the other hand, however, a smaller tolerance will not necessarily increase the accuracy, since on a machine using 16-digit arithmetic the sum will be, at best, good to 16 digits.

The particular series being evaluated is determined by the arguments SOL and FNC within subroutine MATH and communicated to this subroutine via argument DUM.

Output is returned via common: varia-

bles F, DMAX, DLAST, KMAX, and KLAST.

SUM calls on one of the functions of Algorithm 352 (Part B.2.1).

```

C      SUBROUTINE SUM (DUM)
      *****
      INTEGER
      *      K,KLAST,KMAX,L,S
      DOUBLE PRECISION
      *      DLAST,DMAX,DUM,F,T
      COMMON
      *      DUM1(1006),S,L,DUM2(6),F,
      *      DMAX,DLAST,KMAX,KLAST,T
      K      = 0
      F      = DUM(0)
      DMAX   = F
      T      = DABS(F)
      KMAX   = 0
      DO 30 KLAST = 1,L
      DLAST = DUM(KLAST)
      F     = F+DLAST
      IF (T-DABS(DLAST)) 10,10,20
10     DMAX = DLAST
      T     = DABS(DMAX)
      KMAX  = KLAST
20     IF (KLAST.LE.S) GO TO 30
      *     IF (DABS(DLAST)/T.GT.1.D-13)
      *         K = 0
      K     = K+1
      IF (K.EQ.3) GO TO 40
30     CONTINUE
      KLAST = L
40     END
      RETURN
  
```

Algorithm 352 (Part C)
BESSEL (Bessel Functions)²
 (Called by MATH)

Comments The subroutine BESSEL evaluates Bessel functions of the first or second kind, according as the argument SOL = 1 or 2, of orders 0, 1, ..., n and argument u, both of which must be nonnegative. Functions of order zero and one are always evaluated, regardless of the value of n. Results are returned via array JY, with element JY(K) containing the function of order K-1.

It should be noted that for SOL = 2 and u = 0, a large negative constant (-10³⁷) is returned as the function value for all orders and no warning is given.

Different methods of computation are used for J₀(u), J₁(u), Y₀(u), and Y₁(u), depending upon whether u < 8, or not. (See subroutines JOJ1, YOY1, and LUKE for details.) The J_n(u), n = 2, 3, ..., m, are computed by means of a continued fraction (see subroutine JNS), whereas the Y_n(u) for corresponding orders are calculated di-

² This subroutine (together with its subsidiary routines) may be removed in toto, with no changes, and used independently as a Bessel function algorithm. The results are good to 14 significant digits or decimal places, whichever is least accurate, with an error of no more than one unit in the last digit or place.

rectly from the recurrence relation:

$$Y_{n+1}(u) = \frac{2n}{u} Y_n(u) - Y_{n-1}(u)$$

BESSEL calls on the subroutines:

JOJ1—referred to as Algorithm 352 (Part C.1)

YOY1—referred to as Algorithm 352 (Part C.2)

LUKE—referred to as Algorithm 352 (Part C.3)

JNS—referred to as Algorithm 352 (Part C.4)

```

C      SUBROUTINE BESSEL (SOL,U,JY,N)
      *****
      INTEGER
      *      N,NN,SOL
      DOUBLE PRECISION
      *      JY(250),U
      NN = MIN(1,N,249)
      IF (U.EQ.0.DO.AND.SOL.EQ.2)
      *      GO TO 80
      IF (U.GE.8.DO) GO TO 30
      GO TO (10,20), SOL
10     CALL JOJ1 (U,JY)
      GO TO 40
20     CALL YOY1 (U,JY)
      GO TO 40
30     CALL LUKE (U,SOL,JY)
40     IF (N.LT.2) GO TO 100
      GO TO (50,60), SOL
50     CALL JNS (JY,U,NN)
      GO TO 100
C      RECURRENCE FORMULA
60     DO 70 K = 2,NN
      JY(K+1) = 2.DO*
      *      DBLE(FLOAT(K-1))*
      *      JY(K)/U-JY(K-1)
70     CONTINUE
      GO TO 100
80     NN = NN+1
      DO 90 K = 1,NN
      JY(K) = -1.D+37
90     CONTINUE
100    END
      RETURN
  
```

Algorithm 352 (Part C.1)
JOJ1 (First Kind)
 (Called by BESSEL)

Comments The subroutine JOJ1 computes the Bessel functions of the first kind, J₀(x) and J₁(x), for x < 8. This is done by evaluating formula 9.1.10 of [1]. The results are returned via array J.

JOJ1 calls no other subprograms.

```

C      SUBROUTINE JOJ1 (X,J)
      *****
      DOUBLE PRECISION
      *      J(2),T(5),X
      COMMON
      *      DUM(1014),T
      T(1) = X/2.DO
      J(1) = 1.DO
      J(2) = T(1)
      T(2) = -T(1)**2
  
```

```

      T(3) = 1.DO
      T(4) = 1.DO
10     T(4) = T(4)*T(2)/T(3)**2
      J(1) = J(1)+T(4)
      T(5) = T(4)*T(1)/(T(3)+1.DO)
      J(2) = J(2)+T(5)
      IF (DMAX1(DABS(T(4)),DABS(T(5)))
      *      .LT.1.D-15) RETURN
      T(3) = T(3)+1.DO
      GO TO 10
      END
  
```

Algorithm 352 (Part C.2)
YOY1 (Second Kind)
 (Called by BESSEL)

Comments The subroutine YOY1 computes the Bessel functions of the second kind, Y₀(x) and Y₁(x), for x < 8. This is done by evaluating formulas 9.1.13 and 9.1.11 of [1]. The results are returned via array Y. YOY1 calls no other subprograms.

```

C      SUBROUTINE YOY1 (X,Y)
      *****
      DOUBLE PRECISION
      *      T(10),X,Y(2)
      COMMON
      *      DUM(1014),T
      T(1) = X/2.DO
      T(2) = -T(1)**2
      Y(1) = 1.DO
      Y(2) = T(1)
      T(7) = 0.DO
      T(10) = -T(1)
      T(3) = 0.DO
      T(4) = 0.DO
      T(5) = 1.DO
10     T(3) = T(3)+1.DO
      T(4) = T(4)+1.DO/T(3)
      T(5) = T(5)*T(2)/T(3)**2
      Y(1) = Y(1)+T(5)
      T(6) = -T(5)*T(4)
      T(7) = T(7)+T(6)
      T(8) = T(5)*T(1)/(T(3)+1.DO)
      Y(2) = Y(2)+T(8)
      T(9) = -T(8)*(2.DO*T(4)+
      *      1.DO/(T(3)+1.DO))
      T(10) = T(10)+T(9)
      IF (DMAX1(DABS(T(6)),DABS(T(9)))
      *      .GE.1.D-15) GO TO 10
      T(2) = .57721566490153286DO+
      *      DLOG(T(1))
      Y(1) = .63661977236758134DO*
      *      (Y(1)*T(2)+T(7))
      Y(2) = .63661977236758134DO*
      *      (Y(2)*T(2)-1.DO/X)+T(10)/
      *      3.1415926535897932DO
      END
      RETURN
  
```

Algorithm 352 (Part C.3)
LUKE
 (Called by BESSEL)

Comments The subroutine LUKE evaluates Bessel functions of order zero and one, of the first or second kind, according as the argument KIND = 1 or 2, for u ≥ 8. The results are returned via the 2-element array JY.

The Bessel function of the third kind (Hankel function), H_v⁽³⁾(u) = J_v(u) + iY_v(u), can be expressed in terms of the

Chebyshev polynomials, $T_n^*(x)$, as follows:

$$H_\nu^{(1)}(u) = \left(\frac{2}{\pi u}\right)^{\frac{1}{2}} e^{i\left(u - \frac{\nu\pi}{2} - \frac{\pi}{4}\right)} \cdot \sum_{k=0}^{\infty} (\alpha_k^{(\nu)} + i\beta_k^{(\nu)}) T_k^*(R/u). \quad (9)$$

We now define $\alpha_k^{(0)} = A_{k+1}$, $\beta_k^{(0)} = B_{k+1}$, $\alpha_k^{(1)} = C_{k+1}$, $\beta_k^{(1)} = D_{k+1}$, $x = R/u$, and $T_k^*(x) = G_{k+1}(x)$. The recurrence relations for the $G_k(x)$ are as follows:

$$\begin{aligned} G_1(x) &= 1, & G_2(x) &= 2x - 1, \\ G_k(x) &= (4x-2)G_{k-1}(x) - G_{k-2}(x), \\ & & k &\geq 3. \end{aligned}$$

If we let $\nu = 0$ and make other appropriate substitutions in (9), while remembering that $e^{i\theta} = \cos \theta + i \sin \theta$, we can separate the real and imaginary parts and get the following relations:

$$\begin{aligned} J_0(u) &= \left(\frac{2}{\pi u}\right)^{\frac{1}{2}} \\ &\cdot \left[\cos \theta \sum_{k=1}^{\infty} A_k G_k(x) - \sin \theta \sum_{k=1}^{\infty} B_k G_k(x) \right], \\ Y_0(u) &= \left(\frac{2}{\pi u}\right)^{\frac{1}{2}} \\ &\cdot \left[\cos \theta \sum_{k=1}^{\infty} B_k G_k(x) + \sin \theta \sum_{k=1}^{\infty} A_k G_k(x) \right], \end{aligned}$$

where $\theta = u - \pi/4$.

Notice that if $\nu = 1$ in (9), then θ is replaced by $\theta - \pi/2$. Also, $\cos(\theta - \pi/2) = \sin \theta$ and $\sin(\theta - \pi/2) = -\cos \theta$. Therefore, proceeding as before, we get

$$\begin{aligned} J_1(u) &= \left(\frac{2}{\pi u}\right)^{\frac{1}{2}} \\ &\cdot \left[\sin \theta \sum_{k=1}^{\infty} C_k G_k(x) + \cos \theta \sum_{k=1}^{\infty} D_k G_k(x) \right], \\ Y_1(u) &= \left(\frac{2}{\pi u}\right)^{\frac{1}{2}} \\ &\cdot \left[\sin \theta \sum_{k=1}^{\infty} D_k G_k(x) - \cos \theta \sum_{k=1}^{\infty} C_k G_k(x) \right]. \end{aligned}$$

The coefficients A , B , C , and D have been computed for $R = 8$ in eq. (9) and are guaranteed to the number of digits given.

LUKE calls no other subprograms.

```

SUBROUTINE LUKE (U,KIND,JY)
C *****
INTEGER
* K,KIND
DOUBLE PRECISION
* A(19),B(19),CS,C(19),
* D(19),G(3),JY(2),R(2),
* S(2),SN,T,U,X
COMMON
* DUM(1014),R,S,G,X,T,SN,CS
C WARNING - THE FOLLOWING DATA
C STATEMENTS ARE NOT IN ASA
C STANDARD FORTRAN

```

```

DATA A /
* .99959506476867287416D0,
* -.53807956139606913D-3,
* -.13179677123361570D-3,
* .151422497048644D-5,
* .15846861792063D-6,
* -.856069553946D-8,
* -.29572343355D-9,
* .6573556254D-10,
* -.223749703D-11,
* -.44821140D-12,
* .6954827D-13,
* -.151340D-14,
* -.92422D-15,
* .15558D-15,
* -.476D-17,
* -.274D-17,
* .61D-18,
* -.4D-19,
* -.1D-19/

```

```

DATA B /
* -.776935569420532136D-2,
* -.774803230965447670D-2,
* .2536541165430796D-4,
* .394273598399711D-5,
* -.10723498299129D-6,
* -.721389799328D-8,
* .73764602899D-9,
* .150687811D-11,
* -.574589537D-11,
* .45996574D-12,
* .2270323D-13,
* -.887890D-14,
* .74497D-15,
* .5847D-16,
* -.2410D-16,
* .265D-17,
* .13D-18,
* -.10D-18,
* .2D-19/

```

```

DATA C /
* 1.00067753586591346234D0,
* .90100725195908183D-3,
* .22172434918599454D-3,
* -.196575946319104D-5,
* -.20889531143270D-6,
* .1028144350894D-7,
* .37597054789D-9,
* -.7638891358D-10,
* .238734670D-11,
* .51825489D-12,
* -.7693969D-13,
* .144008D-14,
* .103294D-14,
* -.16821D-15,
* .459D-17,
* .302D-17,
* -.65D-18,
* .4D-19,
* .1D-19/

```

```

DATA D /
* .2337682998628580328D-1,
* .2334680122354557533D-1,
* -.3576010590901382D-4,
* .560863149492627D-5,
* .13273894084340D-6,
* .916975845066D-8,
* -.86838880371D-9,
* -.378073005D-11,
* .663145586D-11,
* -.50584390D-12,
* -.2720782D-13,
* .985381D-14,
* -.79398D-15,
* -.6757D-16,
* .2625D-16,
* -.280D-17,
* -.15D-18,
* .10D-18,
* -.2D-19/

```

```

X = 8.0D0/U
G(1) = 1.0D0
G(2) = 2.0D0*X-1.0D0
R(1) = A(1)+A(2)*G(2)
S(1) = B(1)+B(2)*G(2)
R(2) = C(1)+C(2)*G(2)
S(2) = D(1)+D(2)*G(2)
DO 10 K = 3,19
G(3) = (4.0D0*X-2.0D0)*G(2)-G(1)
R(1) = R(1)+A(K)*G(3)
S(1) = S(1)+B(K)*G(3)
R(2) = R(2)+C(K)*G(3)
S(2) = S(2)+D(K)*G(3)
G(1) = G(2)

```

```

G(2) = G(3)
10 CONTINUE
T = .7978845608028654D0/DSQRT(U)
SN = DSIN(U-.7853981633974483D0)
CS = DCOS(U-.7853981633974483D0)
GO TO (20,30), KIND
20 JY(1) = T*(R(1)*CS-S(1)*SN)
JY(2) = T*(R(2)*SN+S(2)*CS)
GO TO 40
30 JY(1) = T*(S(1)*CS+R(1)*SN)
JY(2) = T*(S(2)*SN-R(2)*CS)
40 RETURN
END

```

Algorithm 352 (Part C.4)
JNS

(Called by BESSEL)

Comments The subroutine JNS evaluates Bessel functions of the first kind, of orders $n = 2, 3, \dots, m$, for argument u , given $J_0(u)$ and $J_1(u)$. From the definition $G_n = J_n(u)/J_{n-1}(u)$ and the recurrence relation,

$$J_{n+1}(u) = (2n/u) J_n(u) - J_{n-1}(u),$$

we can derive the following equation:

$$G_n = \frac{1}{\frac{2n}{u} - G_{n+1}}. \quad (10)$$

Since G_{n+1} is of the same form as G_n , we can continue the process and obtain the continued fraction,

$$G_n = \frac{1}{\frac{2n}{u} - \frac{1}{\frac{2(n+1)}{u} - \dots - \frac{1}{\frac{2(n+k)}{u} - G_{n+k+1}}}}. \quad (11)$$

G_m is evaluated using (11), then the other G_n are computed from (10) for $n = m - 1, m - 2, \dots, 2$. Finally, the J_n are evaluated in a forward direction from $J_n = G_n J_{n-1}$ and returned via argument array JJ. See [2] for a more detailed treatment of this process.

JNS calls no other subprograms.

```

SUBROUTINE JNS (JJ,U,M)
C *****
INTEGER
* K,KA,KK,M
DOUBLE PRECISION
* A,B,D(2),DM,G(249),
* JJ(250),P(3),Q(3),U
EQUIVALENCE
* (A,G),(D,G(2)),
* (P,G(4)),(Q,G(7)),
* (DM,G(10)),(B,G(11))
COMMON
* DUM(1014),G,M,K,KK,KA
DM = 2*M
P(1) = 0.0D0
Q(1) = 1.0D0
P(2) = 1.0D0
Q(2) = DM/U
D(1) = P(2)/Q(2)
A = 2.0D0

```

```

10 B = (DM+A)/U
P(3) = B*P(2)-P(1)
Q(3) = B*Q(2)-Q(1)
D(2) = P(3)/Q(3)

IF (DABS(D(1)-D(2))
* .LT.1.D-15) GO TO 20

P(1) = P(2)
P(2) = P(3)
Q(1) = Q(2)
Q(2) = Q(3)
D(1) = D(2)
A = A+2.D0

GO TO 10

20 G(M) = D(2)
KA = M-2
DO 30 K = 1,KA
KK = M-K
A = 2*KK
G(KK) = U/(A-U*G(KK+1))

IF (G(KK).EQ.0.D0)
* G(KK) = 1.D-35

30 CONTINUE
DO 40 K = 2,M
JJ(K+1) = G(K)*JJ(K)
40 CONTINUE

RETURN
END

```

Algorithm 352 (Part B.2.1)
DS, DC, DDS, DDC, PS, PC, DPS, DPC
(Called by MATH via SUM)

Comments The following collection of function subprograms is utilized by SUM to evaluate the k th term ($k = 0, 1, \dots$) of one of the following: eq. (2), (3), (5), (6), or their derivatives.

DS and DC call on functions FJ and FY.

DDS and DDC call on functions FJ, FY, DJ and DY.

PS, PC, DPS, and DPC call no other subprograms.

```

C DOUBLE PRECISION FUNCTION DS(KK)
*****
INTEGER
* K, KK, N, N1, N2, P, S

DOUBLE PRECISION
* AB, FJ, FY

COMMON
* DUM1(1004), N, P, S, DUM2(17),
* K, N1, N2, DUM3(583), AB(200)

C EVALUATES ONE TERM OF THE RADIAL
C SOLUTION, ASSOCIATED WITH B(Q)
K = KK
N1 = K-S
N2 = K+S+P
DS = AB(K+1)*(FJ(N1)*FY(N2)-
* FJ(N2)*FY(N1))

IF (MOD(K+N,2).NE.0) DS = -DS

RETURN
END

C DOUBLE PRECISION FUNCTION DC(KK)
*****
INTEGER
* K, KK, N, N1, N2, P, S

DOUBLE PRECISION
* AB, FJ, FY

COMMON
* DUM1(1004), N, P, S, DUM2(17),
* K, N1, N2, DUM3(583), AB(200)

C EVALUATES ONE TERM OF THE RADIAL
C SOLUTION, ASSOCIATED WITH A(Q)
K = KK

```

```

N1 = K-S
N2 = K+S+P
DC = AB(K+1)*(FJ(N1)*FY(N2)+
* FJ(N2)*FY(N1))

IF (MOD(K+N,2).NE.0) DC = -DC

IF (S+P.EQ.0) DC = .5D0*DC

RETURN
END

C DOUBLE PRECISION FUNCTION DDS(KK)
*****
INTEGER
* K, KK, N, N1, N2, P, S

DOUBLE PRECISION
* AB, DJ, DY, FJ, FY, U1, U2

COMMON
* DUM1(1000), U1, U2, N, P, S,
* DUM2(17), K, N1, N2,
* DUM3(583), AB(200)

C EVALUATES ONE TERM OF THE DERIVATIVE
C OF THE RADIAL SOLUTION,
C ASSOCIATED WITH B(Q)
K = KK
N1 = K-S
N2 = K+S+P
DDS = AB(K+1)*(U2*(FJ(N1)*DY(N2)-
* FJ(N2)*DY(N1))-U1*(FY(N2)*
* DJ(N1)-FY(N1)*DJ(N2)))

IF (MOD(K+N,2).NE.0) DDS = -DDS

RETURN
END

C DOUBLE PRECISION FUNCTION DDC(KK)
*****
INTEGER
* K, KK, N, N1, N2, P, S

DOUBLE PRECISION
* AB, DJ, DY, FJ, FY, U1, U2

COMMON
* DUM1(1000), U1, U2, N, P, S,
* DUM2(17), K, N1, N2,
* DUM3(583), AB(200)

C EVALUATES ONE TERM OF THE DERIVATIVE
C OF THE RADIAL SOLUTION,
C ASSOCIATED WITH A(Q)
K = KK
N1 = K-S
N2 = K+S+P
DDC = AB(K+1)*(U2*(FJ(N1)*DY(N2)+
* FJ(N2)*DY(N1))-U1*(FY(N2)*
* DJ(N1)+FY(N1)*DJ(N2)))

IF (MOD(K+N,2).NE.0) DDC = -DDC

IF (S+P.EQ.0) DDC = .5D0*DDC

RETURN
END

C DOUBLE PRECISION FUNCTION PS(K)
*****
INTEGER
* K, P

DOUBLE PRECISION
* AB, X

COMMON
* DUM1(1005), P, DUM2(2), X,
* DUM3(600), AB(200)

C EVALUATES ONE TERM OF THE ODD
C PERIODIC SOLUTION
PS = AB(K+1)*
* DSIN(DBLE(FLOAT(2*K+P))*X)

RETURN
END

```

```

C DOUBLE PRECISION FUNCTION PC(K)
*****
INTEGER
* K, P

DOUBLE PRECISION
* AB, X

COMMON
* DUM1(1005), P, DUM2(2), X,
* DUM3(600), AB(200)

C EVALUATES ONE TERM OF THE EVEN
C PERIODIC SOLUTION
PC = AB(K+1)*
* DCOS(DBLE(FLOAT(2*K+P))*X)

RETURN
END

C DOUBLE PRECISION FUNCTION DPS(K)
*****
INTEGER
* K, P

DOUBLE PRECISION
* AB, T, X

COMMON
* DUM1(1005), P, DUM2(2), X,
* DUM3(14), T, DUM4(584),
* AB(200)

C EVALUATES ONE TERM OF THE DERIVATIVE
C OF THE ODD PERIODIC SOLUTION
T = 2*K+P
DPS = AB(K+1)*T*DCOS(T*X)

RETURN
END

C DOUBLE PRECISION FUNCTION DPC(K)
*****
INTEGER
* K, P

DOUBLE PRECISION
* AB, T, X

COMMON
* DUM1(1005), P, DUM2(2), X,
* DUM3(14), T, DUM4(584),
* AB(200)

C EVALUATES ONE TERM OF THE DERIVATIVE
C OF THE EVEN PERIODIC SOLUTION
T = 2*K+P
DPC = -AB(K+1)*T*DSIN(T*X)

RETURN
END

```

Algorithm 352 (Part B.2.2)
FJ, FY, DJ, DY (Bessel Functions and Derivatives)
(Called by DS, DC, DDS, DDC)

Comments The following collection of function subprograms produces Bessel functions or their derivatives for integer order n , n being positive or negative. This is accomplished by using the already computed functions of nonnegative order (Algorithm 352 (Part C)) and substituting them in one of the following formulas:

$$J_{-n}(u) = (-1)^n J_n(u),$$

$$Y_{-n}(u) = (-1)^n Y_n(u),$$

$$J_n'(u) = \frac{n}{u} J_n(u) - J_{n+1}(u),$$

$$Y_n'(u) = Y_{n-1}(u) - \frac{n}{u} Y_n(u),$$

whichever is appropriate.
 DJ calls on function FJ.
 DY calls on function FY.
 FJ and FY call no other subprograms.

```

C   DOUBLE PRECISION FUNCTION FJ(N)
C   *****
C   INTEGER
C   *       K,N
C   DOUBLE PRECISION
C   *       J
C   COMMON
C   *       J(250),DUM(527),K
C   PRODUCES BESSEL FUNCTIONS
C   OF THE FIRST KIND
C   K = IABS(N)
C   IF (K.GE.250) GO TO 20
C   FJ = J(K+1)
C   IF (MOD(N,2).LT.0) FJ = -FJ
10   RETURN
20 FJ = 0.D0
   WRITE (6,99) N
99  FORMAT(2HOJ13,7H NEEDED)
   GO TO 10
C   END
C   DOUBLE PRECISION FUNCTION FY(N)
C   *****

```

```

INTEGER
*       K,N
DOUBLE PRECISION
*       Y
COMMON
*       DUM1(500),Y(250),DUM2(27),K
C   PRODUCES BESSEL FUNCTIONS
C   OF THE SECOND KIND
C   K = IABS(N)
C   IF (K.GE.250) GO TO 20
C   FY = Y(K+1)
C   IF (MOD(N,2).LT.0) FY = -FY
10   RETURN
20 FY = 0.D0
   WRITE (6,99) N
99  FORMAT(2HOY13,7H NEEDED)
   GO TO 10
C   END
C   DOUBLE PRECISION FUNCTION DJ(N)
C   *****
C   INTEGER
C   *       N
C   DOUBLE PRECISION
C   *       FJ,FN,U1
C   COMMON
C   *       DUM1(1000),U1,DUM2(26),FN
C   DERIVATIVES OF BESSEL FUNCTIONS
C   OF THE FIRST KIND

```

```

FN = N
IF (N-249) 10,20,40
10 DJ = FN*FJ(N)/U1-FJ(N+1)
   GO TO 30
20 DJ = FJ(N-1)-FN*FJ(N)/U1
   GO TO 30
30   RETURN
40 DJ = 0.D0
   WRITE (6,99) N
99  FORMAT(3HOJ@13,7H NEEDED)
   GO TO 30
C   END
C   DOUBLE PRECISION FUNCTION DY(N)
C   *****
C   INTEGER
C   *       N
C   DOUBLE PRECISION
C   *       FN,FY,U2
C   COMMON
C   *       DUM1(1002),U2,DUM2(24),FN
C   DERIVATIVES OF BESSEL FUNCTIONS
C   OF THE SECOND KIND
C   IF (N.GE.250) GO TO 20
C   FN = N
C   DY = FY(N-1)-FN*FY(N)/U2
10   RETURN
20 DY = 0.D0
   WRITE (6,99) N
99  FORMAT(3HOY@13,7H NEEDED)
   GO TO 10
C   END

```

REMARK ON ALGORITHM 268 [R2]
 ALGOL 60 REFERENCE LANGUAGE EDITOR
 [W. M. McKeeman, *Comm. ACM* 8 (Nov. 1965), 667]
 G. SAUER (Recd. 23 Dec. 1968)
 Institut für Theoretische Physik der Justus-Liebig-Universität, 63 Giessen, West Germany
 KEY WORDS AND PHRASES: symbol manipulation
 CR CATEGORIES: 4.49

In the *procedure send*, replace the line
 1 until 1 do if *buffer*[*u*+1] =
 with the line
 1 until *tabstop* do if *buffer*[*u*+1] = (1)
 The published version fails to clear the buffer when a line to be printed contains no blanks and *tabstop* > 0, causing an array bounds violation. Knowing *buffer*[*tabstop*+1] never to contain a blank character, the search for blanks may be stopped at *u* = *tabstop* + 1.

(1) The author is indebted to the referee for suggesting this brief form.

The policy concerning the contributions of algorithms to *Communications of the ACM* appears, most recently, in the January 1969 issue, page 39. A contribution should be in the form of an algorithm, a certification, or a remark. An algorithm must normally be written in the ALGOL 60 Reference Language or in USASI Standard FORTRAN or Basic FORTRAN.

REMARK ON ALGORITHM 274 [F1]
 GENERATION OF HILBERT DERIVED TEST MATRIX [J. Boothroyd, *Comm. ACM* 9 (Jan. 1966), 11]
 J. BOOTHROYD (Recd. 7 Jan. 1969)
 University of Tasmania, Hobart, Tasmania, Australia
 KEY WORDS AND PHRASES: test matrix, Hilbert matrix
 CR CATEGORIES: 5.14

An alternative, simpler, and more efficient procedure for generating test matrices having the same properties as those generated by Algorithm 274 is given below. The method, like that of Algorithm 274, is due to T. J. Dekker and may be described as follows.

The elements of the inverse of a segment of a Hilbert matrix are given by

$$(H^{-1}) = (-1)^{i+j} \times f_i \times f_j / (i + j - 1)$$

where

$$f_i = \text{factorial}(n + i - 1) / (\text{factorial}(i - 1) \uparrow 2 / \text{factorial}(n - i)).$$

The f_i may be factored as $f_i = f_{i1} \times f_{i2}$, in which

$$f_{i1} = \binom{n + i - 1}{i - 1} \times n, \quad f_{i2} = \binom{n - 1}{n - i}.$$

Test matrices T are constructed by $T = D_1 H D_2$ where $D_1 = \text{diag}(f_{i1})$, $D_2 = \text{diag}(f_{i2})$, and H is the Hilbert matrix segment $H_{i,j} = 1/(i + j - 1)$. It may be seen that this is equivalent to defining the T matrices by:

$$T_{i,j} = (f_{i1} f_{j1}) / (i + j - 1),$$

$$f_{i1} = \binom{n + i - 1}{i - 1} \times n, \quad f_{j1} = \binom{n - 1}{n - j},$$

with fi, fj given by the recurrence relations:

$$(fi)_1 = n, \quad (fi)_{i+1} = (fi)_i \times (n + i)/i,$$

$$(fj)_1 = 1, \quad (fj)_{j+1} = (fj)_j \times (n - j)/j.$$

That the condition $K(T)$ of these matrices is severe may be seen from an observation of the referee, who notes that

$$K(T) = \|T\| \times \|T^{-1}\|,$$

$$\geq (\max_{1 \leq i, j \leq n} t_{i,j})^2 = (t_{n, (n+1)} \div 2)^2 \sim (2 \uparrow 3n / 13n)^2,$$

where $\|\cdot\|$ is the L_1, L_2, L_∞ , or the Euclidean matrix norm.

Other properties of these matrices shared by those of Algorithm 274 are:

- (a) Each matrix has unit determinant;
- (b) The eigenvalues form a set $\lambda_1, \lambda_2, \dots, 1/\lambda_2, 1/\lambda_1$, so that odd order matrices have one eigenvalue of unity.

The procedure *testmx1* below has been tested on an Elliott 503 (positive integer word length of 38 bits) and matrices of all orders up to 13 were generated before integer overflow occurred with $n = 14$.

procedure *testmx1* (*a, n*); **value** *n*; **integer** *n*; **array** *a*;
comment generates in $a[1 : n, 1 : n]$ test matrices with integer elements given by

$$t_{i,j} = \binom{n+i-1}{i-1} \times n \times \binom{n-1}{n-j} / (i+j-1)$$

and such that the elements of T inverse are $(-1)^{i+j} \times t_{i,j}$.

To determine for a particular computer that limit on n which permits the exact machine representation of all elements of these matrices, the following maximum values are listed:

<i>n</i>	$t_{i,j}$ (max)
8	163800
9	1178100
10	8314020
11	61108047
12	440936496;

begin

```

integer i, j, fi, fj, illess1;
fi := n; illess1 := 0;
for i := 1 step 1 until n do
  begin
    fj := 1;
    for j := 1 step 1 until n do
      begin
         $a[i, j] := (fi \times fj) \div (illess1 + j);$ 
         $fj := ((n-j) \times fj) \div j$ 
      end;
       $fi := ((n+i) \times fi) \div i; illess1 := i$ 
    end
  end testmx1

```

Proofs that the test matrices described above have integer elements and checkerboard inverses follow the lines of similar proofs given in [1].

Acknowledgments: Thanks are due to T. J. Dekker for communicating details of this method and to the referee for the contribution mentioned.

REFERENCE:

1. DEKKER, T. J. Evaluation of determinants, solution of systems of linear equations and matrix inversion. Rep. No. MR63, Mathematical Centre, Amsterdam, June 1963, pp. 8 and 9.

REMARK ON ALGORITHM 333 [H]
 MINIT ALGORITHM FOR LINEAR PROGRAMMING [Rodolfo C. Salazar and Subrata K. Sen, *Comm. ACM* 11 (June 1968), 437]

D. K. MESSHAM (Recd. 27 Nov. 1968 and 28 Feb. 1969)
 Nelson Research Laboratories, The English Electric Co. Ltd., Stafford, England

KEY WORDS AND PHRASES: linear programming, dual simplex method, primal problem, dual problem
 CR CATEGORIES: 5.41

The procedure has been tested with Marconi Myriad Algol, and it ran successfully when the following changes had been made (the first is merely a misprint):

1. The first statement in procedure *results* was changed from $z := e[1 : lcol];$ to $z := e[1, lcol];$
2. To satisfy an ALGOL 60 restriction that a type procedure should contain an assignment to its procedure identifier, the **real** on the first line of the procedure was removed.

3. It is possible for the published algorithm to give incorrect results when it reaches a state in *phase1* where there are no possible pivotal elements in one column of the tableau. (For example, maximize $-x_1 - x_2 - x_3$, with $2x_1 + x_2 = 3$ and $x_3 = 1$, reaches this state.) To correct this the line in procedure *phase1*

```

if gamma < gmin then
  was changed to
if gamma < gmin  $\wedge$  thmin [ind[k]] < 106 then
  All the appearances of 106 in this algorithm should be written as 106.

```

The following improvements are also suggested:

4. It is assumed that *lcol* is a global integer with the correct value. This was made unnecessary by adding *lcol* to the list of integers declared on the line immediately following the initial comment; the bounds of the array *ind*, declared on the next line, were changed

```

  from [1 : lcol]
  to [1 : m+n-p+1];
  and lcol := m + n - p + 1;

```

was inserted as the first executable statement of the procedure *MINIT* (after **end phase1**);

5. It is assumed that equality constraints will be given with positive right-hand sides. This restriction was overcome by inserting in the procedure *phase1* after the line **integer array imin** [1 : *lcol*]; the following:

```

for i := m - p + 2 step 1 until m + 1 do
if  $e[i, lcol] < 0$  then
for j := 1 step 1 until lcol do  $e[i, j] := -e[i, j];$ 

```

PLAN TO ATTEND

ACM 69

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