

7. Summary

We have introduced a formalism which allows us to explicate certain rather gross properties of language processing systems. As it is, the notation should be useful for designing the outlines of complex programming systems and their implementation, and it should be especially good for documentation. The formalism should also provide a mathematical basis which can be extended to handle more detailed properties of such systems. Some specific inadequacies where it could be extended follow.

1. It does not describe the amount of compilation or interpretation, unless it is coupled with precise definitions of the languages involved. For instance, in (7) we have no idea whether IL is close to machine language or to the source language. IL could be little more than assembly language, or just a trivial modification of the source language, or anything in between. Of course precise definitions of SL, IL, and ML would clear this up.

2. It does not permit the description of such processes as incremental compilation.

3. It does not permit the formal description of systems involving programs which consist of two or more pieces written in different languages, such as FSL.

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Algorithms

L. D. FOSDICK, Editor

ALGORITHM 395

STUDENT'S *t*-DISTRIBUTION [S14]

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KEY WORDS AND PHRASES: Student's *t*-statistic, distribution function, approximation, asymptotic expansion

CR CATEGORIES: 5.12, 5.5

real procedure *student* (*t*, *n*, *normal*, *error*); **value** *t*, *n*; **real** *t*, *n*;
real procedure *normal*, *error*;

comment *student* evaluates the two-tail probability $P(t | n)$ that *t* is exceeded in magnitude for Student's [1] *t*-distribution with *n* degrees of freedom. The procedure provides results accurate to 11 decimal places and 8 significant digits for integer values of *n*, with approximate continuation of the function through noninteger values of *n* (over 6 decimal places for $n > 4.3$).

The procedure *normal* (*x*) returns the area under the standard normal frequency curve to the left of *x*, so that a negative argument yields the lower-tail area. The user-supplied procedure, *error*(*n*), should produce a diagnostic warning and may go to a label, terminate, or return a distinctive value (zero or -1.0) as a signal of error to the calling program.

Student's series expansion of the probability integral is supplemented by a faster asymptotic approximation for large values of *n* and by a more precise "tail" series expansion for large values of *t*.

The value of *x*, defined as the normal deviate at the same probability level as *t*, may be approximated by an asymptotic normalizing expansion of Cornish-Fisher type [2].

$$x = z + (z^3 + 3z)/b - (4z^7 + 33z^5 + 240z^3 + 855z)/10b^2 \\ + (64z^{11} + 788z^9 + 9801z^7 + 89775z^5 + 543375z^3 + 1788885z)/210b^3 - \dots$$

where $z = (a \times \ln(1 + t^2/n))^{1/2}$, $a = n - \frac{1}{2}$ and $b = 48a^2$ [3].

This is well approximated by the first three terms with the third term's divisor replaced by

$$10b(b + 0.8z^4 + 100).$$

The *student* probability is double the normal single-tail area, corresponding to the deviate *x*.

The maximum error in the probability result for all values of *t* is displayed as a function of *n* in Figure 1, for this approximation, for the first few terms of the asymptotic expansion and for Fisher's [4] fifth-order approximation used in Algorithm 321 [5] for $n \geq 30$.

For small *n* and moderate *t* the result is calculated as $P(t | n) = 1 - A(t | n)$ using Student's cosine series for $A(t | n)$, rearranging formulas 26.7.3 and 26.7.4 of the NBS Handbook [6] in nested form

$$A(t | n \text{ odd}) = \frac{2}{\pi} \left[a_1 \tan(y) + \frac{y}{b} \left\{ 1 + \frac{2}{3b} \left\{ \dots \frac{(n-5)}{(n-4)b} \right. \right. \right. \\ \left. \left. \left. \cdot \left\{ 1 + \frac{(n-3)}{(n-2)b} \right\} \dots \right\} \right\} \right] \\ A(t | n \text{ even}) = \frac{y}{\sqrt{b}} \left\{ 1 + \frac{1}{2b} \left\{ \dots \frac{(n-5)}{(n-4)b} \left\{ 1 + \frac{(n-3)}{(n-2)b} \right\} \dots \right\} \right\},$$

where $y = \sqrt{(t^2/n)}$ and $b = 1 + t^2/n$. In the nested form, terms are treated in reverse order to the summation in Algorithm 321 and Algorithm 344 [7], reducing the number of operations required and reducing build up of roundoff error. Explicit decre-

menting of the "loop" parameter ensures that its final value remains defined on exit from the loop for use in an odd/even test.

Execution times for Fortran versions run on a CDC 3200 with programmed floating point are displayed in Figure 2, which indicates that nesting decreases the time for the cosine series method by about 30 percent and that it is appropriate to change over to the asymptotic method (using Algorithm 209 [8] for *normal*) when $n \geq 20$. Although this approximation would be accurate to more than 11 decimal places, the use of Algorithm 209 limits accuracy to about 9 decimals. This accuracy may be sufficient for many applications, in which case *student* may be abbreviated by deleting lines 15 and 27 through 35, removing the declaration and assignment of z from line 3, replacing line 5 by

if $n > \text{entier}(n) \vee n \geq 20$ then

and replacing line 25 by

$\text{student} := \text{if } a > 1.0 \text{ then } 0.0 \text{ else } 1.0 - a$

The latter avoids spurious negative results due to roundoff error when a is near 1 for large values of t . The storage required for this abbreviated version was a little less than for Algorithm 344 and less than half that for Algorithm 321.

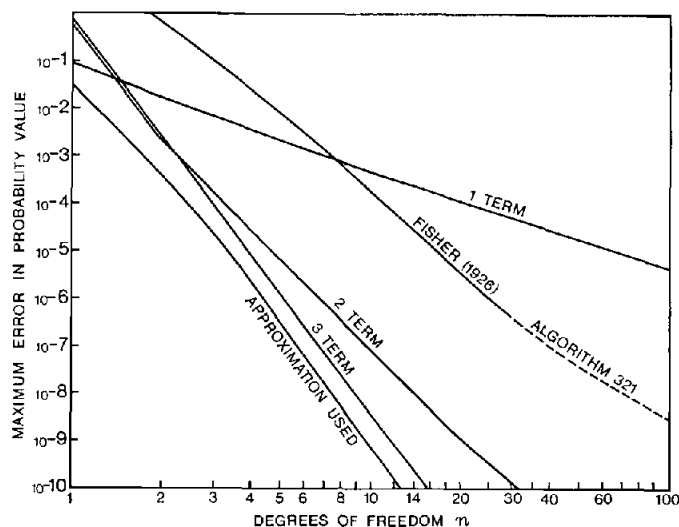


FIG. 1. Maximum error of approximations for "Student's" t -probability: 1, 2, and 3 term expansion, approximation with adjusted divisor, and Fisher's 5th order approximation

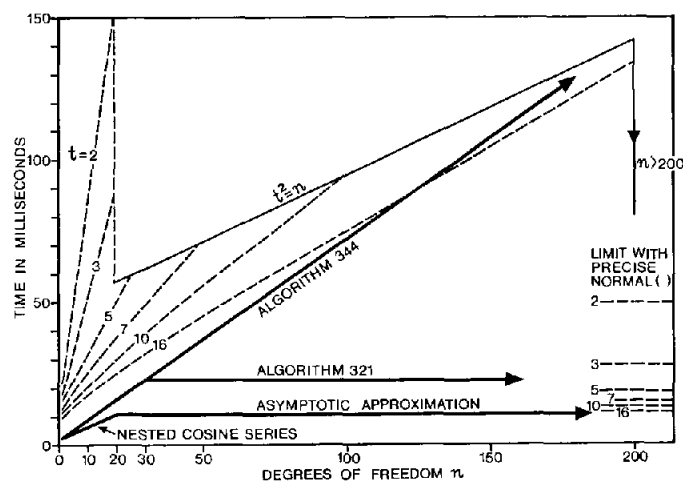


FIG. 2. Execution times (CDC3200 with programmed floating point). Broken lines: "tail" series for selected values of t (upper left); asymptotic method using precise *normal* (right)

Applications such as production of tables or function inversion to obtain extreme quantiles may require greater precision at extreme probability levels than these methods provide. For the cosine series and the asymptotic approximation using a high precision procedure for *normal*, such as Algorithm 304 [9], the relative error in the result increases in magnitude as the result decreases to extremely small values, as illustrated in Figure 3.

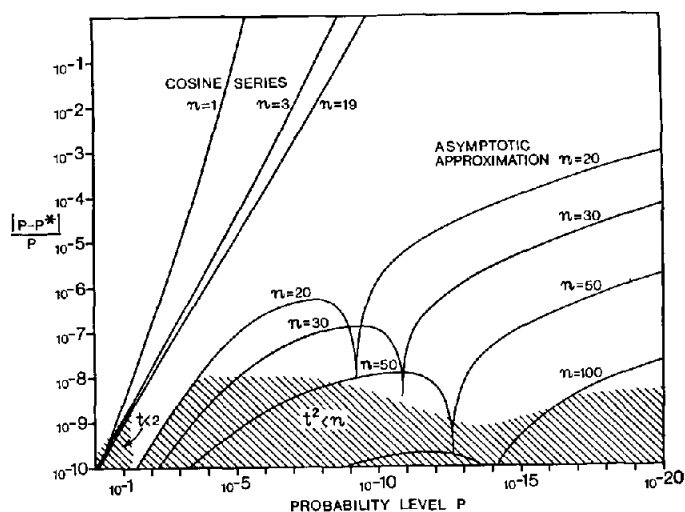


FIG. 3. Relative error, $|P - P^*|/P$, of approximation P^* ; shaded region for restricted t values

For small P more precise results are obtained using a series expansion of $P(t | n)$ in terms of $w = 1/\sqrt{1+t^2/n}$,

$$P(t|n) = C(n) \times w^n \left\{ \frac{1}{n} + \frac{1 \times w^2}{2(n+2)} + \frac{1 \times 3 \times w^4}{2 \times 4(n+4)} + \dots \right\},$$

where $C(n) = \Gamma((n+1)/2)/(\sqrt{\pi} \times \Gamma(n/2))$. The series is summed till a negligible term occurs and then the factor $C(n) \times w^n$ is applied using the same repeated loop as the cosine series. Except for w near 1 when t is small, the truncation error is small, and accumulation of error in the repeated loop is moderate unless n is very large.

The cosine series method loses precision mainly in the subtraction $1 - A(t | n)$ as well as from the *sqrt* procedure and *arctan* when n is odd. In the worst case, $n = 19$, the error is kept below 3 decimals by changing to the tail series if $t > 2$, which ensures 8 significant digits in the result for the 36-bit (about 11 decimal) precision real variables for the processor used. As shown in Figure 3, change over from the asymptotic method to the tail series when $t^2 > n$ maintains about 8 significant digits in the result. For a machine of greater precision the use of more terms in the asymptotic series may be warranted, and the change over criteria would need adjustment to balance speeds and precision between the three methods.

Execution times for the tail series are shown as broken lines in Figure 2 for selected values of t : with bounds $t \geq 2$ for $n < 20$, $t^2 \leq n$ for $n \geq 20$ and with the limit $n < 200$ preventing excessive time for large t beyond a probability level near 10^{-10} . For the asymptotic method, using for *normal* a higher precision procedure based on Algorithm 304, the execution times for different values of the argument approach those shown at the right of Figure 2. Averaged over a range of arguments arising in practice, the provision for higher precision more than doubles the time required. In the case of Smirnov's [10] 6D tables of $S(t | n) = 1 - 0.5 \times P(t | n)$, retabulation to 10D, using the more precise procedure for *normal*, increased the time from about 7 minutes to 12 minutes, while introducing the tail series method to tabulate $P(t | n)$ over the same range to 8 significant digits increased the time further to about 16 minutes. Use of the asymptotic

approximation enabled Smirnov's 6D tables of $\psi(t | 1000/\xi)$, which is an approximate continuation of $S(t | n)$ over non-integer values of $n = 1000/\xi$, to be extended to 10D for $\xi = 0(2)30$ in 5 minutes, and permits continuation to $\xi = 200$ with over 6D accuracy as indicated in Figure 1.

The preparation of diagrams by Murray C. Childs is gratefully acknowledged.

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if $n < 1$ then student := error(n) else

begin

real a, b, y, z ; $z := 1.0$;

$t := t \uparrow 2$; $y := t/n$; $b := 1.0 + y$;

if $n > \text{entier}(n) \vee n \geq 20 \wedge t < n \vee n > 200$ then

begin

comment Asymptotic series for large or noninteger n ;

if $y > n-6$ then $y := \ln(b)$;

$a := n - 0.5$; $b := 48.0 \times a \uparrow 2$; $y := a \times y$;

$y := ((((-0.4 \times y - 3.3) \times y - 24.0) \times y - 85.5) /$

$(0.8 \times y \uparrow 2 + 100.0 + b) + y + 3.0) / b + 1.0) \times \text{sqr}(y)$;

student := $2.0 \times \text{normal}(-y)$;

end

else

if $n < 20 \wedge t < 4.0$ then

begin

comment Nested summation of "cosine" series;

$a := y := \text{sqr}(y)$; if $n = 1$ then $a := 0.0$;

loop:

$n := n - 2$; if $n > 1$ then

begin $a := (n-1)/(b \times n) \times a + y$; go to loop end;

$a :=$ if $n = 0$ then $a/\text{sqr}(b)$

else $(\arctan(y) + a/b) \times 0.63661977236$;

comment $2/\pi = 0.6366197723675813430755351 \dots$;

student := $z - a$

end

else

begin

comment "tail" series expansion for large t -values;

integer j ; $a := \text{sqr}(b)$; $y := a \times n$; $j := 0$;

for $j := j + 2$ while $a \neq z$ do

begin

$z := a$; $y := y \times (j-1)/(b \times j)$; $a := a + y/(n+j)$

end;

$n := n + 2$; $z := y := 0.0$; $a := -a$; go to loop

end

end

ALGORITHM 396

STUDENT'S t -QUANTILES [S14]

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KEY WORDS AND PHRASES: Student's t -statistic, quantile, asymptotic approximation

CR CATEGORIES: 5.12, 5.5

real procedure t quantile ($P, n, \text{normdev}, \text{error}$);

value P, n ; real P, n ; real procedure $\text{normdev}, \text{error}$;

comment This algorithm evaluates the positive quantile at the (two-tail) probability level P , for Student's t -distribution with n degrees of freedom. The quantile function is an inverse of the two-tail

$$P(t|n) = 2 \frac{\Gamma(\frac{1}{2}n + \frac{1}{2})}{\sqrt{(\pi n)\Gamma(\frac{1}{2}n)}} \int_t^\infty \frac{du}{(1+u^2/n)^{(n+1/2)}}$$

which is approximated in Algorithm 395 [1] by series whose inverses are used in this algorithm for t quantiles. Test calculations to 36-bit precision indicate that the result is correct to at least 6 significant digits, even for the analytic continuation through noninteger values of $n > 5$.

The procedure $\text{normdev}(p)$ is assumed to return a negative normal deviate at the lower tail probability level p , e.g. -2.32 for $p = 0.01$. The user-supplied procedure for $\text{error}(n)$ should give a diagnostic warning that the value of P or n is invalid and may go to a label, terminate, or return a distinctive value as an error signal to the calling program.

For $n = 1$ and $n = 2$ the exact result of integration is readily inverted to yield $t = \cot(P \times \pi/2)$ and $t^2 = 2/(P(2-P)) - 2$, respectively. For larger n an asymptotic inverse expansion about normal deviates is applicable, while for smaller values of P a second series expansion is used to achieve sufficient precision. Both approximations have been adjusted to enhance precision for n as low as 3.

Both methods involve an expansion of the factor

$$d/n = \frac{1}{2} \sqrt{\pi} \Gamma(\frac{1}{2}n) / \Gamma(\frac{1}{2}n + \frac{1}{2})$$

in terms of $a = 1/(n - \frac{1}{2})$ and $b = 48/a^2$

$$d/n = \sqrt{(a\pi/2)} (1 - 3/b + 94.5/b^2 - 9058.5/b^3 + \dots) [2].$$

A three term approximation uses $b(b+c)$ instead of b^2 as a divisor, where the coefficients in

$$c = 96.36 - 16a - 98a^2 + 20700a^3/b,$$

have been fitted to ensure 8 significant digits in d for n as low as 3.

The inverse asymptotic expansion of Cornish-Fisher type relates a function $y(t) = \sqrt{[(n - \frac{1}{2}) \ln(1 + t^2/n)]}$ to the normal deviate x at the corresponding probability level, $P/2$:

$$y = x - (x^3 + 3x)/b + (4x^7 + 63x^5 + 360x^3 + 945x)/10b^2$$

$$- (64x^{11} + 1628x^9 + 19881x^7 + 145719x^5 + 694575x^3$$

$$+ 1902285x)/210b^3 + \dots [2],$$

whence $t = \sqrt{[n \times (\exp(a \times y^2) - 1)]}$. For a three term approximation the third term's divisor is replaced by

$$10b \times (b + c - 2x - 7x^2 - 5x^3 + 0.05 \times d \times x^4),$$

whose coefficients have been fitted to reduce the error for small n and for larger n and x . For $n < 5$, c is increased by $0.3(n - 4.5)(x + 0.6)$ to further reduce error in an interval of P not well covered by the following approximation.

For small P , where t^2/n is large, the integrand may be ex-

panded in terms of $w^2 = 1/(1+t^2/n)$ and integrated term by term to yield

$$P = \frac{nw^n}{d} \left\{ \frac{1}{n} + \frac{w^2}{2(n+2)} + \frac{1 \times 3w^4}{2 \times 4(n+4)} + \dots \right\},$$

which may be inverted to express t^2/n in terms of $y = (P \times d)^{2/n}$

$$\frac{t^2}{n} = \frac{1}{y} + \frac{n+1}{n+2} \left\{ -1 + \frac{y}{2(n+4)} + \frac{n \times y^2}{3(n+2)(n+6)} + \frac{n(n+3)(2n^2+9n-2)y^3}{8(n+2)^2(n+4)^2(n+8)} + \dots \right\}.$$

Since the ratio of successive terms is nearly $n \times y/(n+6)$ for small n , replacement of the term in y^2 by $y/[3(n+2)\{(n+6)/(n \times y) - 1.0\}]$ provides an approximate allowance for subsequent terms in the series, which is empirically improved by replacing the -1.0 by $-0.822 - 0.089 \times d$.

As n and P increase, the errors for the asymptotic approximation decrease, whereas errors for the second series increase, so that for each value of n the error curves intersect at a value of P above which the asymptotic approximation is better and below which the second series should be used. By adjusting the two approximations the error level at these intersections has been balanced at about the seventh significant digit for $n \geq 3$ and $P > 10^{-24}$. The value of y at these points is about $a + 0.05$ and this fact provides a convenient criterion for selecting which approximation to use: the asymptotic series if y exceeds $a + 0.05$, otherwise the second series.

Although better approximations could be obtained by use of more terms in each series, greater precision can be achieved by using the result of this algorithm as a starting value for iterative inversion of $P(t | n)$, whose value and derivative can be computed with considerable precision using recurrence relations as in Algorithm 395.

A comparison of results from this algorithm against values obtained by inverting the function provided by Algorithm 395 indicates a precision of over 6 significant digits for $10^{-24} \leq P \leq 0.9$, $n \geq 1$. At the conventional tabulation points in $0.001 \leq P \leq 0.9$ results for $n = 1$, $n = 2$, and $n > 10$ checked to 8 significant digits.

Previously published tables [3, 4, 5] provide 3 or 4 decimal place check values, some of which are found to be slightly in error. Thus for $n = 2$, $P = 0.001$, t is given as 31.598 by Fisher and Yates and by Federighi, 31.5991 by Smirnov, and 31.5990546 by this procedure, while for $n = 1$, $P = 0.001$ the value 636.6096 given by Smirnov conflicts with Fisher and Yates, Federighi (636.619) and this procedure (636.61925). Other errors in the last few digits in Smirnov's table for low values of n and P include 10.2129 for $n = 3$, $P = 0.002$, which should be 10.2145, and 4.7812 for $n = 9$, $P = 0.001$, which should be 4.7809.

t quantile may be used to obtain percentiles at values of P and n not provided in existing tables or for extending their accuracy. Such tables are customarily used for assessing the significance of a sample value for t , but for automatic computation the probability level is more effectively determined as $P(t | n)$ using a direct procedure such as Algorithm 395.

Pseudorandom t -values may be generated for sampling applications by using uniformly distributed pseudorandom numbers for P , and in this case *normdev* may be a **real procedure** returning pseudorandom normal deviates which are independent of P .

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if $n < 1 \vee P > 1.0 \vee P \leq 0.0$ then t quantile := error(n)
else if $n = 2$ then t quantile := $\text{sqrt}(2.0/(P \times (2.0 - P)) - 2.0)$

else

begin

real half pi; half pi := 1.5707963268;

if $n = 1$ then

begin $P := P \times \text{half pi}$; t quantile := $\cos(P)/\sin(P)$ end

else

begin

real a, b, c, d, x, y ;

$a := 1.0/(n - 0.5)$; $b := 48.0/a \uparrow 2$;

$c := ((20700 \times a/b - 98) \times a - 16) \times a + 96.36$;

$d := ((94.5/(b+c) - 3.0)/b + 1.0) \times \text{sqrt}(a \times \text{half pi}) \times n$;

$x := d \times P$; $y := x \uparrow (2.0/n)$;

if $y > 0.05 + a$ then

begin

comment Asymptotic inverse expansion about normal;

$x := \text{normdev}(P \times 0.5)$; $y := x \uparrow 2$;

if $n < 5$ then $c := c + 0.3 \times (n - 4.5) \times (x + 0.6)$;

$c := (((0.05 \times d \times x - 5.0) \times x - 7.0) \times x - 2.0) \times x + b + c$;

$y := (((((0.4 \times y + 6.3) \times y + 36.0) \times y + 94.5)/c - y - 3.0)/b + 1.0) \times x$;

$y := a \times y \uparrow 2$;

$y := \text{if } y > 0.002 \text{ then } \exp(y) - 1.0 \text{ else } 0.5 \times y \uparrow 2 + y$

end

else $y := ((1.0/(((n+6.0)/(n \times y) - 0.089 \times d - 0.822) \times (n+2.0) \times 3.0) + 0.5/(n+4.0)) \times y - 1.0) \times (n+1.0)/(n+2.0) + 1.0/y$;

t quantile := $\text{sqrt}(n \times y)$

end

end Student's t -quantile

ALGORITHM 397

AN INTEGER PROGRAMMING PROBLEM [H]

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KEY WORDS AND PHRASES: integer programming, change-making problem

CR CATEGORIES: 5.41

procedure MINDIST($C, M, \text{SENSE}, W, \text{RESULT}$);

value C, M ; integer C, M ; Boolean SENSE ;

integer array W, RESULT ;

comment This algorithm solves an integer programming problem described in [1]. Given is a fixed weight vector $w = (w_1, w_2, \dots, w_m)$, where the w_i are nonnegative integers, where m is a positive integer, and where

$$1 = w_1 < w_2 < \dots < w_m$$

For any nonnegative integer c (representing cost), an m -distribution of c relative to w is an m -tuple $(a_1, a_2, \dots, a_m)$ such that the a_i are nonnegative integers, and such that $\sum_{i=1}^m a_i w_i = c$. The m -distribution $(a_1, a_2, \dots, a_m)$ is minimal if, for any m -distribution $(b_1, b_2, \dots, b_m)$ of c relative to w , we have $\sum_{i=1}^m a_i \leq \sum_{i=1}^m b_i$. The m -distribution $(a_1, a_2, \dots, a_m)$ is standard if it is obtainable as follows:

$$\begin{aligned} c_m &= c \\ c_i &= c_{i+1} - a_{i+1} \times w_{i+1} \quad (i=m-1, m-2, \dots, 1) \\ a_i &= c_i / w_i \quad (i=m, m-1, \dots, 1) \end{aligned}$$

(where all divisions are integer divisions).

If $MINDIST(C, M, SENSE, W, RESULT)$ is called with a nonnegative integer C , a positive integer M , and an array $W = (W[1], W[2], \dots, W[M])$, then the resulting array

$RESULT = (RESULT[1], RESULT[2], \dots, RESULT[M])$ is a minimal M -distribution of C relative to W . If, before calling $MINDIST$, $SENSE$ is set to **true**, then $MINDIST$ retains $SENSE$ as **true** if and only if $RESULT$ is also a standard M -distribution of C relative to W .

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```
begin
  integer I, J, R, Q, SUM, SUN;
  integer array A[1:M], B[1:M];
  if M = 1 then
    begin
      RESULT[1] := C;
    EXIT1 :
      go to EXIT
    end
  Q := C/W[M];
  if (Q×W[M]) > C then Q := Q - 1;
  R := C - W[M] × Q;
  if M = 2 then
    begin
      RESULT[1] := R; RESULT[2] := Q;
    EXIT2 :
      go to EXIT
    end;
  J := 0;
LOOP:
  MINDIST (R+J×W[M], M-1, SENSE, W, B);
  if J ≠ 0 then go to NOT ZERO;
BETA:
  for I := 1 step 1 until M-1 do A[I] := B[I];
  A[M] := 0;
GAMMA:
  if J = Q then
    begin
      for I := 1 step 1 until M do RESULT[I] := A[I];
    EXIT3:
      go to EXIT
    end;
  SUM := 0;
  for I := 1 step 1 until M do SUM := SUM + A[I];
  if (W[M]×SUM-R-J×W[M])/(W[M]-W[M-1]) ≤ 0 then
    begin
      for I := 1 step 1 until M-1 do RESULT[I] := A[I];
      RESULT[M] := A[M] + Q - J;
    EXIT4:
      go to EXIT
    end;
  J := J + 1;
  go to LOOP;
```

NOT ZERO:

```
SUM := 0; SUN := 0;
for I := 1 step 1 until M do SUM := SUM + A[I];
for I := 1 step 1 until M-1 do SUN := SUN + B[I];
if SUM ≤ SUN then
  begin A[M] := A[M] + 1; go to GAMMA end;
SENSE := false;
go to BETA;
EXIT:
end PROCEDURE MINDIST
```

ALGORITHM 398

TABLELESS DATE CONVERSION* [Z]

RICHARD A. STONE (Recd. 2 Jan. 1970 and 6 April 1970)

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Princeton, NJ 08540

* Patent applied for.

KEY WORDS AND PHRASES: date, calendar

CR CATEGORIES: 5.9

```
procedure calendar(y, n, m, d);
  value y, n; integer y, n, m, d, t;
comment calendar is called with the year in y and the day of the
year in n. The month number is returned in m, and the day of the
month is returned in d. The first section of the procedure changes
the dates so that February has 30 days. The second section uses
the fact that 30.55 (m+2) - 91 passes through the number of
days preceeding each month.
  Error detection: m will be in the range 1-12 if and only if n
is in the correct range;
begin
  t := if (y ÷ 4) * 4 = y then 1 else 0;
comment The following statement is unnecessary
  if it is known that 1900 < y < 2100;
  t := if (y ÷ 400) * 400 = y ∨ (y ÷ 100) * 100 ≠ y then t else 0;
  d := n + (if n > (59+t) then 2 - t else 0);
  m := ((d+91)*100) ÷ 3055;
  d := (d+91) - (m*3055) ÷ 100;
  m := m - 2
end calendar
```

ALGORITHM 399

SPANNING TREE [H]

JOUKO J. SEPPÄNEN (Recd. 6 Jan. 1970 and 8 May 1970)

Computing Center, Helsinki University of Technology,

Otaniemi, Finland

KEY WORDS AND PHRASES: graph, tree, spanning tree

CR CATEGORIES: 5.32

```
procedure spanning tree(v, e, I, J, p, T);
  value v, e; integer v, e, p; integer array I, J, T;
```

comment This procedure grows a spanning tree T for a given undirected loop-free graph $G = (N, E)$ of v vertices and e edges. If G is disconnected a spanning forest will be grown.

The edges $(I[k], J[k]) \in E$ for $k = 1, 2, \dots, e$ are assumed to be stored in the arrays $I[1:e]$ and $J[1:e]$. At each stage of the algorithm one edge is considered whereby one of four possible conditions will arise. If neither of the vertices is included in a tree, this edge is taken as a new tree and its vertices numbered by an incremented component number c . If one vertex is in a tree, the edge will be grown to this tree. If the two vertices are in different trees, these will be grafted into a single tree by renumbering the vertices of the other component. Finally, if both vertices are in the same tree, the edge completes a fundamental cycle of the graph with respect to the spanning tree and consequently will not be considered further. At the end, the indices of the edges in the spanning tree are stored in the array $T[1:v-p]$ where p is the number of trees in the forest. The procedure can also be used to find a minimal spanning tree by sorting the edges into ascending order before calling the procedure.

The main loop in the procedure is executed e times. For cases where the ratio e/v is high it could be worthwhile to introduce an additional variable, say d , in the program, for keeping a count of the number of edges included in T . When d has attained the value of $v - 1$ the algorithm could terminate.

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5. KNUTH, D. E. *The Art of Computer Programming, Vol I Fundamental Algorithms*. Addison-Wesley, Reading, Mass., 1968. pp. 370-371;

begin

```

integer i, j, k, c, n, r;
integer array V[1:v];
c := n := 0;
for k := 1 step 1 until v do V[k] := 0;
for k := 1 step 1 until e do
begin
i := I[k]; j := J[k];
if V[i] = 0 then
begin
T[k-n] := k;
if V[j] = 0 then V[i] := V[j] := c := c + 1
else
V[i] := V[j]
end
end
else if V[j] = 0 then
begin
T[k-n] := k; V[j] := V[i]
end
else if V[i] ≠ V[j] then
begin
T[k-n] := k; i := V[i]; j := V[j];
for r := 1 step 1 until v do
if V[r] = j then V[r] := i
end graft
else n := n + 1
end edge;
p := v - e + n
end spanning tree

```

ALGORITHM 400

MODIFIED HAVIE INTEGRATION [D1]

GEORGE C. WALLICK (Recd. 26 Jan. 1970 and 25 Apr. 1970)

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KEY WORDS AND PHRASES: numerical integration, Havie integration, Romberg quadrature, modified Romberg-quadrature, trapezoid values, rectangle values

CR CATEGORIES: 5.16

DESCRIPTION:

The Havie integration method for the approximate evaluation of the definite integral

$$I = \int_A^B F(x) dx \quad (1)$$

as implemented in ACM Algorithm 257 [4] is based upon the parallel generation of the Romberg table of trapezoidal T_j^k values [1] and the table of rectangular R_j^k values also used by Krasun and Prager [3]. At each step in the development of the tables the difference $|T_j^k - R_j^k|$ is examined. If $|T_j^k - R_j^k| \leq \epsilon$ the process is said to have converged and the algorithm returns a value of

$$T_j^{k+1} = \frac{1}{2}(T_j^k + R_j^k). \quad (2)$$

For some $F(X)$, e.g. $F(X) = e^{-X^2}$ and $F(X) = 2/(2 + \sin 10\pi X)$, the R_j^k , T_j^k pairs converge more rapidly than the Romberg sequence of T_j^k values. (This is the same class of $F(X)$ for which a simple nonadaptive Simpsons Rule algorithm [5] is competitive with the Havie algorithm.) For other $F(X)$, the Havie algorithm is slightly less efficient than the Romberg algorithm.

Like Romberg quadrature, Havie integration requires the evaluation of the rectangular values

$$R_0^k = \frac{B-A}{2^k} \sum_{j=1}^{2^k} F \left[A + (j-\frac{1}{2}) \frac{B-A}{2^k} \right]. \quad (3)$$

Rutishauser [6] recognized that this repeated addition of small terms to a large partial sum can lead to serious roundoff error. He suggested a procedure for the evaluation of the R_0^k which significantly reduces this error. The method, used by Fairweather [2] in a modified Romberg algorithm, leads to a significant improvement in accuracy for large orders of extrapolation.

In the modified Havie integration algorithm HRVINT the R_0^k are evaluated using a 3-level version of the Rutishauser procedure. The arguments X of the generating function $F(X)$ are evaluated as in eq. (3) rather than by accumulative addition as in Algorithm 257.

In the argument list for HRVINT, F is the name of the generating function FUNCTION $F(X)$ which returns a value of $F(X)$ corresponding to a specified value of X , A , and B represent the lower and upper limits of integration, and MAX is the maximum order of extrapolation to be permitted, $\text{MAX} \leq 16$. Values of $\text{MAX} > 16$ are interpreted as $\text{MAX} = 16$; the value of MAX is not changed by the subprogram. Computation is terminated when

$$|T_j^k - R_j^k| \leq \text{ACC} * |T_j^k|$$

or when the order of extrapolation $\text{MFIN} = \text{MAX}$. Here ACC is a measure of the desired relative accuracy, $\text{ACC} > 0$. Upon exit HRVINT is the approximate value of the integral, FAC is a measure of the final relative accuracy achieved

$$\text{FAC} = |T_j^k - R_j^k| / |T_j^k|$$

and MFIN is the order of extrapolation.

Test case. HRVINT was tested in Fortran IV on a CDC 6400 computer using single-precision floating point arithmetic (14+ decimal digits). Corresponding integral values were also obtained

using a Fortran version of the standard Havie Algorithm 257. The results of these tests are summarized in Table I.

For modest accuracy requirements, the two algorithms are seen to be equivalent. For both algorithms the maximum accuracy achievable is limited by truncation and roundoff error. Since the Rutishauser modification serves to reduce the magnitude of such errors, the modified Havie algorithm can, in many cases, return optimum integral values that are from 1 to 2 significant figures more accurate than those returned by Algorithm 257.

In the routine use of the algorithms it is possible to specify an

accuracy requirement that cannot be satisfied. When this condition obtains, the algorithms are forced to proceed to the maximum permitted extrapolation order. With Algorithm 257 error accumulation accompanying such an overspecification can lead to a serious decline in evaluation accuracy. With the modified Havie algorithm HRVINT this loss is minimized and in most cases virtually eliminated.

Acknowledgment. The author wishes to thank Mobil Research and Development Corporation for permission to publish this information.

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ALGORITHM:

```

FUNCTION HRVINT(F,A,B,MAX,ACC,FAC,MFIN)
C HAVIE INTEGRATION WITH AN EXPANDED RUTISHAUSER-
C TYPE SUMMATION PROCEDURE
  DIMENSION T(17),U(17),TPREV(17),UPREV(17)
C TEST FOR MAX GREATER THAN 16
  MUX=MAX
  IF(MAX-16)10,10,5
  MUX=16
C INITIALIZATION
  10 ENPT=0.5*(F(A)+F(B))
  SUMT=0.0
  MFIN=1
  N=1
  H=B-A
  SH=H
C BEGIN REPETITIVE LOOP FROM ORDER 1 TO ORDER MAX
  15 T(1)=H*(ENPT+SUMT)
  SUM=0.
  NN=N+N
  EN=NN
  EM=SH/EN
C BEGIN RUTISHAUSER EVALUATION OF RECTANGULAR SUMS
C INITIALIZATION
  IF(NN-16)20,20,25
  20 NZ=NN
  GO TO 30
  25 NZ=16
  IF(NN-256)30,30,35
  30 NA=NN
  GO TO 40
  35 NA=256
  IF(NN-4096)40,40,45
  40 NB=NN
  GO TO 50
  45 NB=4096
C DEVELOPMENT OF RECTANGULAR SUMS
  50 DD 70 KC=1,NN,4096
  SUMB=0.
  KK=KC+NB-1
  DD 65 KB=KC,KK,256
  SUMA=0.
  KKK=KB+NA-1
  DD 60 KA=KB,KKK,16
  SUMZ=0.
  KFR=KA+NZ-1
  DD 55 KZ=KA,KFR,2
  ZKZ=KZ
  SUMZ=SUMZ+F(A+ZKZ*EM)
  55 SUMA=SUMZ+SUMA
  60 SUMB=SUMA+SUMB
  65 SUM=SUMB+SUM
C END OF RUTISHAUSER PROCEDURE
  U(1)=H*SUM
  K=1
C BEGIN EXTRAPOLATION LOOP
  75 FAC=ABS(T(K)-U(K))
  IF(T(K))80,85,80
C TEST FOR RELATIVE ACCURACY
  80 IF(FAC-ABS(ACC*T(K)))90,90,100
C TEST FOR ABSOLUTE ACCURACY WHEN T(K)=0
  85 IF(FAC-ABS(ACC))95,95,100
  90 FAC=FAC/ABS(T(K))
C INTEGRAL EVALUATION BEFORE EXIT
  95 HRVINT=0.5*(T(K)+U(K))
  RETURN
  100 IF(K-MFIN)105,115,115
  105 AK=K+K
  D=2.**AK
  DMA=D-1.0

```

TABLE I. A COMPARISON OF THE HAVIE AND MODIFIED HAVIE ALGORITHMS

$$I = \int_A^B F(X) dX$$

(m = Extrapolation Order, $m \leq 16$; N.S.F. = Number of Significant Figures)

F(X)	A	B	Correct value (digits 10-16)	Numerical Evaluation						
				Specified relative accuracy	Havie			Modified Havie		
					I (digits 10-14)	m	N.S.F.	I (digits 10-14)	m	N.S.F.
e^{-x^2}	0.0	5.0	45139 55	10^{-1} - 10^{-2} 10^{-2} - 10^{-10} 10^{-11} 10^{-12} 10^{-13} 10^{-14} 10^{-16}	46726 45039 45110 45128 45134 39757 39757	3 4 5 6 6 16 16	10 11 12 12 12 9 9	46726 45039 45111 45131 45137 45137 45136	3 4 5 6 6 7 10	10 11 12 12 13 13 13
$\ln x$	1.0	10.0	29940 46	10^{-9} 10^{-10} 10^{-11} - 10^{-12} 10^{-13} 10^{-14}	29845 29937 29937 29937 29556	8 8 9 9 16	11 13 13 13 11	29846 29939 29940 29940 29940	8 8 9 10 10	11 13 14 14 14
$(1+x)^{-1}$	0.0	1.0	55994 53	10^{-9} 10^{-10} 10^{-11} 10^{-12} 10^{-13} 10^{-14} - 10^{-16}	56353 55996 55990 55988 55987 53242	6 6 6 7 8 16	11 13 13 12 12 10	56354 55997 55991 55991 55991 55991	6 6 6 7 7 9	11 13 13 13 13 13
$(1+x^4)^{-1}$	0.0	1.0	33991 10	10^{-6} - 10^{-7} 10^{-8} - 10^{-10} 10^{-11} - 10^{-12} 10^{-13} 10^{-14} - 10^{-16}	35633 33993 33984 30854 30854	5 6 7 16 16	10 13 12 10 10	35634 33995 33989 33987 33988	5 6 7 7 9	10 13 13 13 13
x^{-2}	0.01	1.1	68595 04	10^{-8} 10^{-9} 10^{-10} 10^{-11} 10^{-12} - 10^{-13} 10^{-14} - 10^{-16}	71022 68136 68076 64508 64508 64508	13 13 13 16 16 16	10 11 10 10 10 10	71529 68647 68589 68590 68589 68584	13 13 13 14 14 16	10 11 12 12 12 12
x^{-4}	0.01	1.1	89506 64	10^{-8} 10^{-9} 10^{-10} 10^{-11} - 10^{-12} 10^{-13} 10^{-14} - 10^{-16}	89368 89199 88857 86878 86878 86878	13 13 14 16 16 16	11 11 10 10 10 10	89694 89526 89503 89502 89502 89499	13 13 14 14 15 16	11 12 13 13 13 12
x^{-8}	0.01	1.1	29246 64	10^{-8} 10^{-9} - 10^{-10} 10^{-11} 10^{-12} - 10^{-13} 10^{-14} 10^{-16}	29556 28828 27557 27557 27557 27557	13 14 16 16 16 16	11 11 10 10 10 10	29767 29247 29245 29244 29244 29242	13 13 14 15 16 16	10 14 13 13 13 13

```

C BEGIN EXTRAPOLATION
  T(K+1)=(D*T(K)-TPREV(K))/DMA
  TPREV(K)=T(K)
  U(K+1)=(D*U(K)-UPREV(K))/DMA
  UPREV(K)=U(K)
C END EXTRAPOLATION
  K=K+1
  IF(K-MUX)75,110,110
C END EXTRAPOLATION LOOP
110  FAC=ABS(T(K)-U(K))
  IF(T(K))90,95,90
C ORDER IS INCREASED BY ONE
115  H=0.5*H
  SUMT=SUMT+SUM
  TPREV(K)=T(K)
  UPREV(K)=U(K)
  MFIN=MFIN+1
  N=NN
  GO TO 15
C RETURN FOR NEXT ORDER EXTRAPOLATION
END

```

REMARK ON ALGORITHM 304[S15]
 NORMAL CURVE INTEGRAL [I. D. Hill and S. A. Joyce, *Comm. ACM* 10(June 1967), 374]
 Bo HOLMGREN (Recd. 30 Apr. 1970)
 Dept. KDO, ASEA, S-721 83 Västerås, Sweden

KEY WORDS AND PHRASES: normal curve integral, probability, special functions
 CR CATEGORIES: 5.12, 5.5

Algorithm 304 with the remark of Adams was translated into Fortran IV and run on a GE-625 computer. The GE-625 has a 28-bit mantissa and allows exponents up to 10^{38} . With *upper* = false and $x < -2.32$, the routine ran into overflow at several values of x . To avoid this the following lines

```

if q2 > 1030 then
begin
  p1 := p1 × 10-30; p2 := p2 × 10-30;
  q1 := q1 × 10-30; q2 := q2 × 10-30
end;

```

were inserted after the line

```

s := m; m := t;

```

REMARK ON ALGORITHM 347 [M1]
 AN EFFICIENT ALGORITHM FOR SORTING WITH
 MINIMAL STORAGE [Richard C. Singleton, *Comm. ACM* 12 (Mar. 1969), 185]
 RICHARD PETO (Recd. 18 Feb. 1970)
 Medical Research Council, 115 Gower Street, London
 W. C. 1

KEY WORDS AND PHRASES: sorting, ranking, minimal storage sorting, digital computer sorting
 CR CATEGORIES: 5.31

If the values of ij , instead of always being $(i+j) \div 2$, are at varying positions between i and j , then there is less likelihood of peculiar initial structure causing failure of the algorithm to perform rapidly. The position of ij can be made to vary by replacing the statements

```

m := 0; ii := i; go to L4; L1: ij := (i+j) ÷ 2;
by
real r; r := 0.375; m := 0; ii := i; go to L4;
L1: r := if r > 0.58984375 then r - 0.21875 else r + 0.0390625;
ij := i + (j-i) × r;

```

comment These four decimal constants, which are respectively 48/128, 75.5/128, 28/128, and 5/128, are rather arbitrary. On most compilers their binary representations will be exact, and the use of them in the statement *L1* causes r to vary cyclically over the 33 values 48/128 ... 80/128. Therefore ij takes a variable position somewhere within the middle quarter of the segment to be sorted. Wider variation of ij would be undesirable in the special case of a partially presorted array;

In sorting an array of N elements which are initially in random order this will waste (on ICL Atlas) less than $N/10^6$ seconds, but if the array is, for example, composed initially of two equal presorted halves, then the use of the original rather than the modified version would more than double the sorting time required if $N > 10^4$.

As the author points out, the published version could fail if used to sort arrays of 1024 or more elements because the upper bounds of *IU* and *IL* might be inadequate. For a standard procedure the declaration *IL, IU* [0:8] should be replaced by the declaration *IL, IU* [0:20]. This permits the sorting of arrays of up to 4 million elements, which is, with present core store sizes, sufficient.

The statement $tt := a[L]$ which precedes *L3*: will be executed less frequently if it is transferred into the next conditional statement, which then reads

```

if k ≤ L then begin tt := a[L]; a[L] := a[k]; a[k] := tt;
go to L2 end

```

REMARK ON ALGORITHM 368 [D5]
 NUMERICAL INVERSION OF LAPLACE
 TRANSFORMS [Harald Stehfest, *Comm. ACM* 13 (Jan. 1970), 47]
 HARALD STEHFEST (Recd. 6 May 1970)
 Institut f. angew. Physik, J. W. Goethe-Universität
 6000 Frankfurt a.M., W. Germany

KEY WORDS AND PHRASES: Laplace transform inversion, integral transformations, integral equations
 CR CATEGORIES: 5.15, 5.18

Some errors have crept into the comment of the procedure after proof-reading:
 The formula following "and thus" should read

$$\sum_{i=1}^K x_i(K) F_{N/2+1-i} = F\left(\frac{\ln 2}{\alpha}\right) + (-1)^{K+1} \alpha^K \frac{(N/2-K)!}{(N/2)!} + o\left(\frac{(N/2-K)!}{(N/2)!}\right).$$

The formula following "with" should read

$$V_i = (-1)^{N/2+i} \sum_{k=\lceil \frac{i+1}{2} \rceil}^{Min(i, N/2)} \frac{k^{N/2} (2k)!}{(N/2-k)! k! (k-1)! (i-k)! (2k-i)!}.$$

The policy concerning the contributions of algorithms to *Communications of the ACM* has been revised and was published in the August 1970 issue, page 513. Copies of "Algorithm Policy / Revised August 1970" will be mailed upon request. Sept 1970 p. 573